

Towards Fine-Grained Secure Communications

The slide features several decorative circles in a light purple color. There are two solid circles in the top row, one solid circle in the bottom row, and two hollow circles in the top row. The text is centered and overlaid on these circles.

Tatsuaki Okamoto
(NTT)

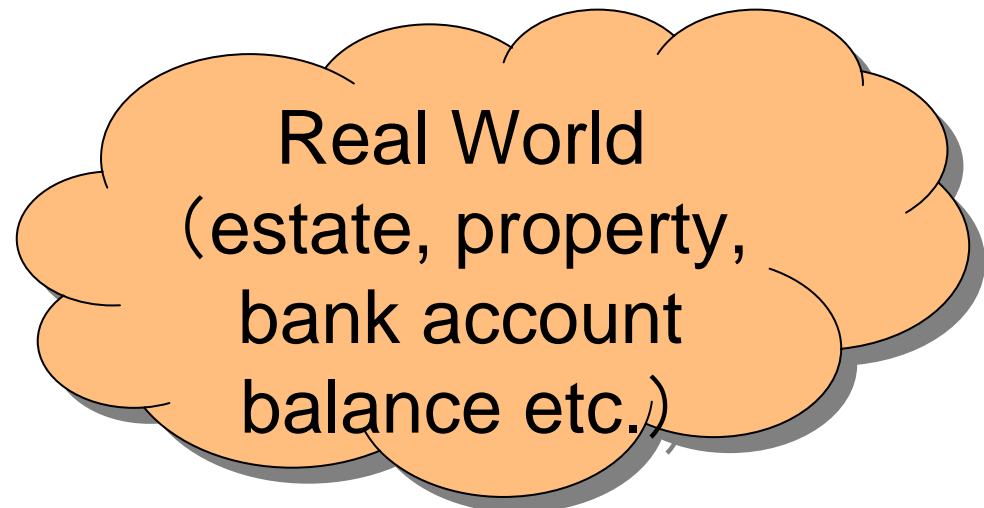
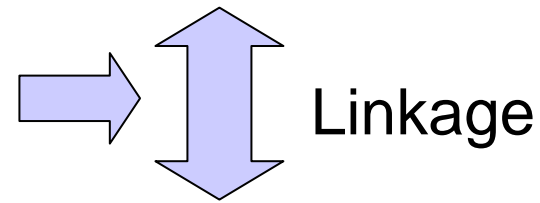
The Role of Cryptography and Information Security

ICT systems . . .

(Every information is communicated and recorded by digital data)

Virtual World

Cryptography and Information Security



Basic Properties of Cryptography

- **Confidentiality**: Data flow control
(a data item is only accessed by qualified persons) ... **encryption**
- **Authentication**: Correctness of data and users/parties ... **signatures, identification**
- **Privacy**: Anonymity of users/parties
... **privacy-enhanced primitives**



Confidentiality

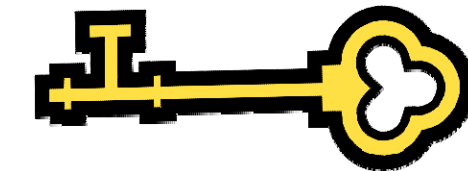
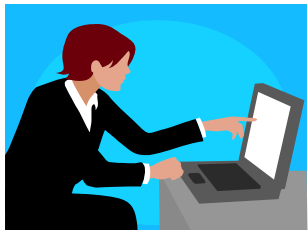
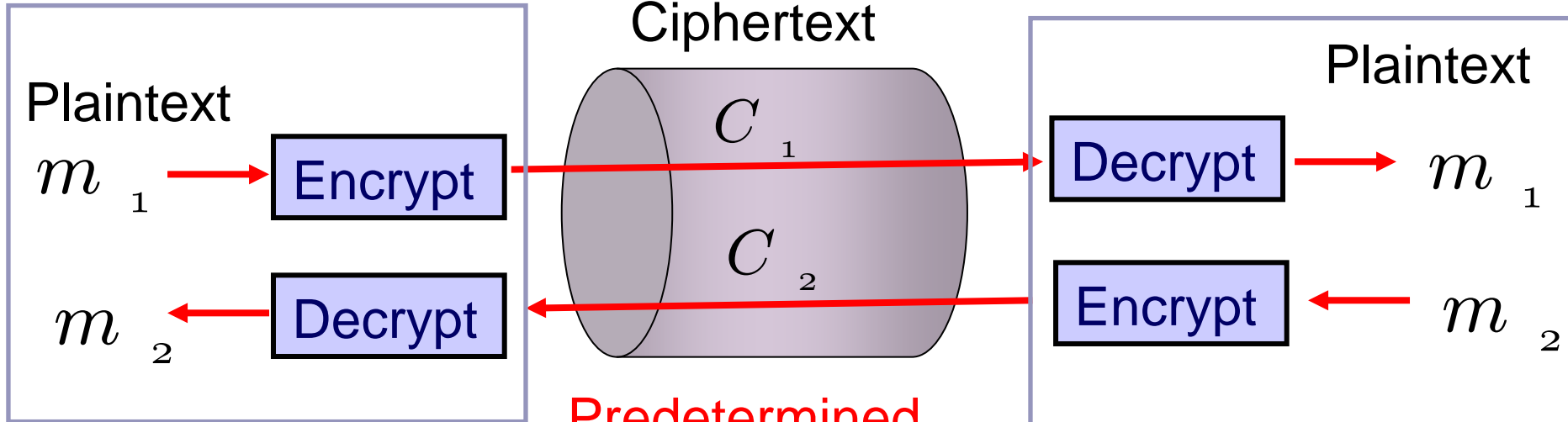
The Most Traditional Way to Realize Confidentiality

- **Symmetric Encryption**: realized by sharing a secret algorithm or a secret key between a pair of sender and receiver.
- **Long history for several thousand years** (similar to that of civilizations of human beings).

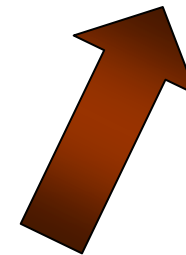
Symmetric Encryption

Alice

Bob



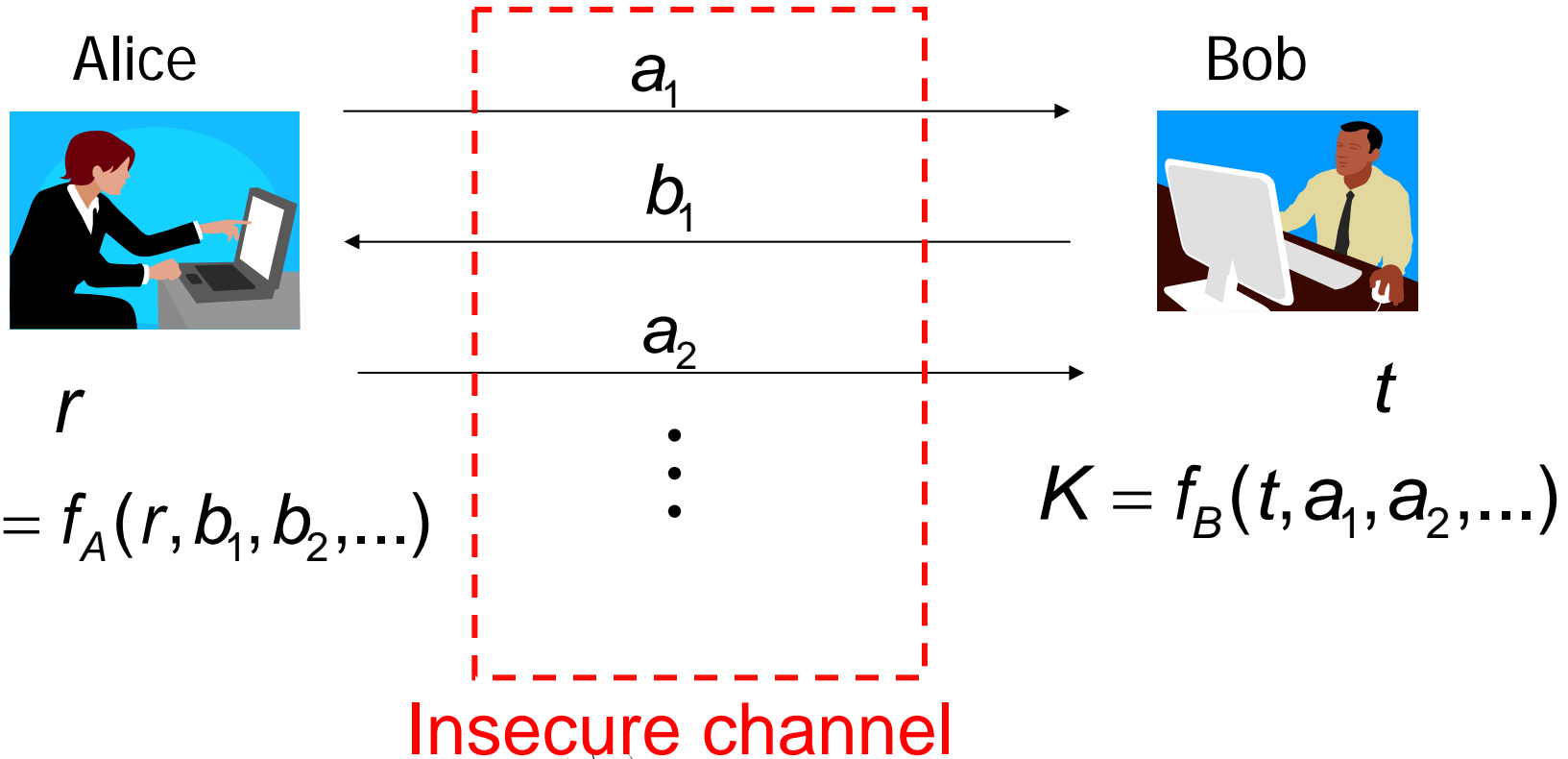
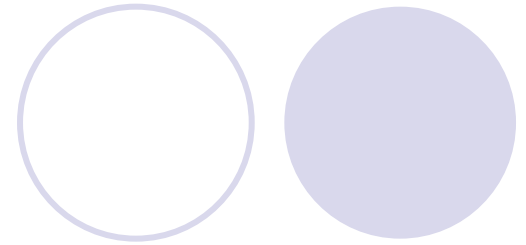
Secret Key /
Secret Algorithm



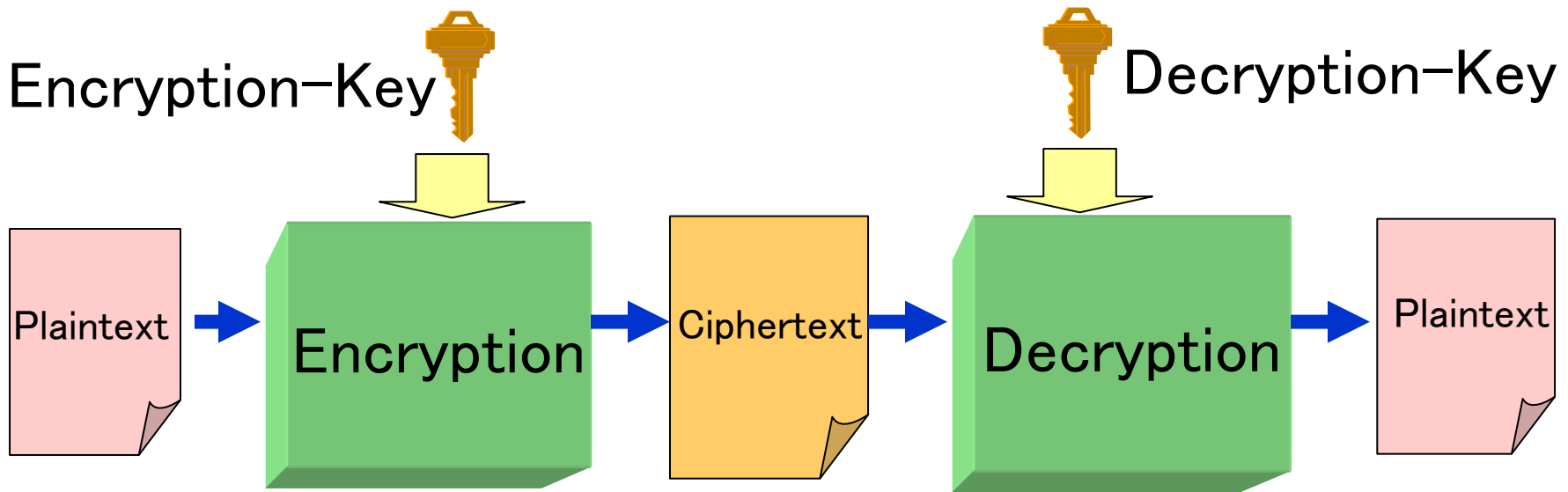
A problem raised in the early 1970's

- **The early 1970's:** ARPANET (the origin of the Internet) started growing.
- **Problem:** Is it possible for two parties to share a secret key securely through an insecure channel, where all messages exchanged by the two parties are wiretapped by the adversary?

Secret Key Sharing over Insecure Channels

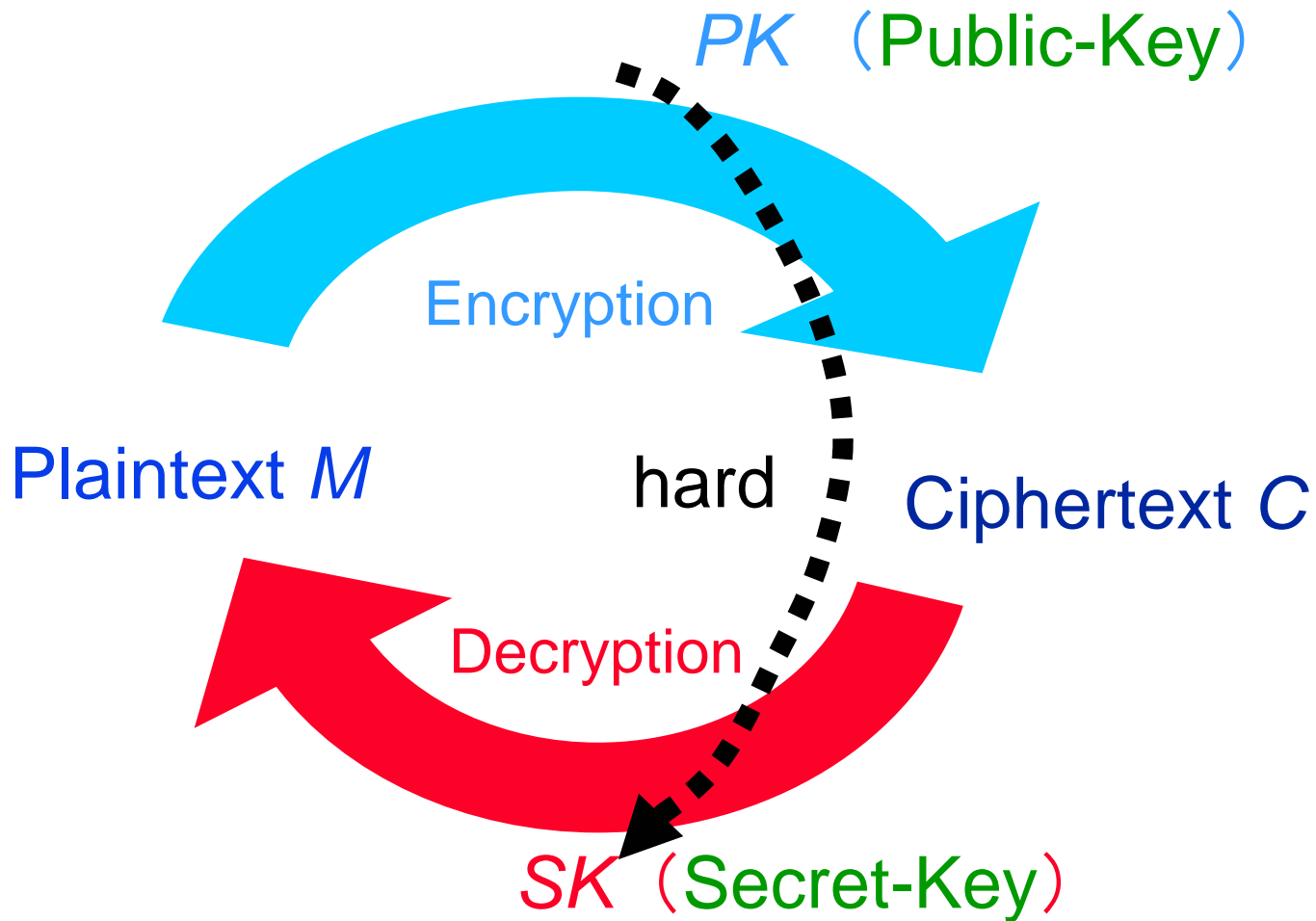


Public-key Encryption (1976)

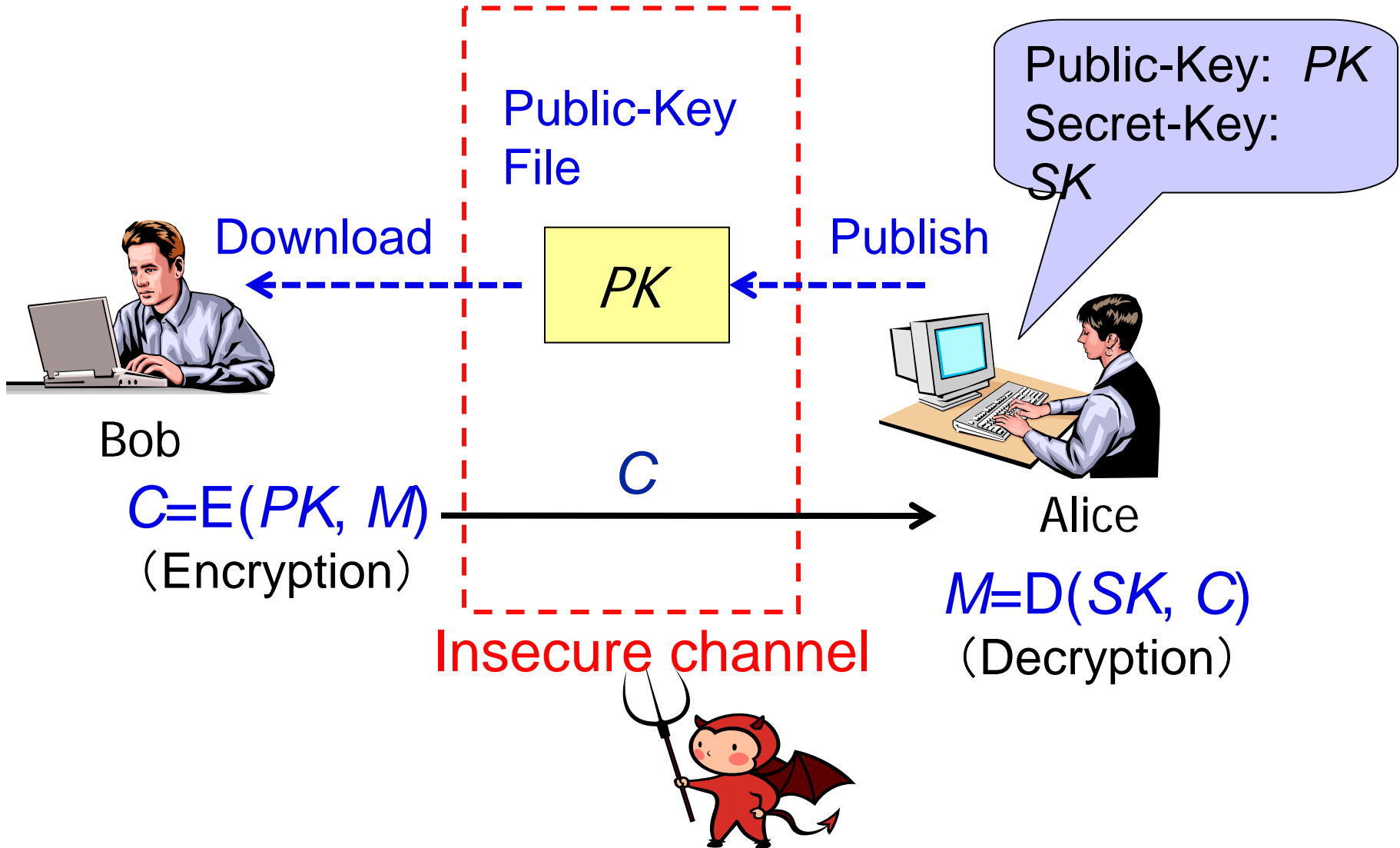


| | |
|-----------------------|--------------------------------------|
| Symmetric Encryption | Encryption-Key = Decryption-Key |
| Public-key Encryption | Encryption-Key \neq Decryption-Key |

Principle of Public-key Encryption



Public-key Encryption



A New Concept of Encryption

- Public-key encryption as well as symmetric encryption are now widely used in many applications, especially in the internet.
- To get a public-key of each receiver may be a problem in some applications.
- Can we use an identity (e.g., email address) of a receiver in place of the public-key?
- No, it is impossible, in the concept of public-key encryption.
- It may be possible in a new concept of encryption.

Identity-Based Encryption

Public-Key
(system parameters)

PK

Publish



Authority

Public-Key: PK
Secret-Key:
 SK
(Master- SK)

Download



Bob



Identity
(e.g., Alice@ibe.com)



Secret-Key
for Alice
(SK_{Alice})



Alice

$$C = E(\text{Alice@ibe.com}, PK, M)$$

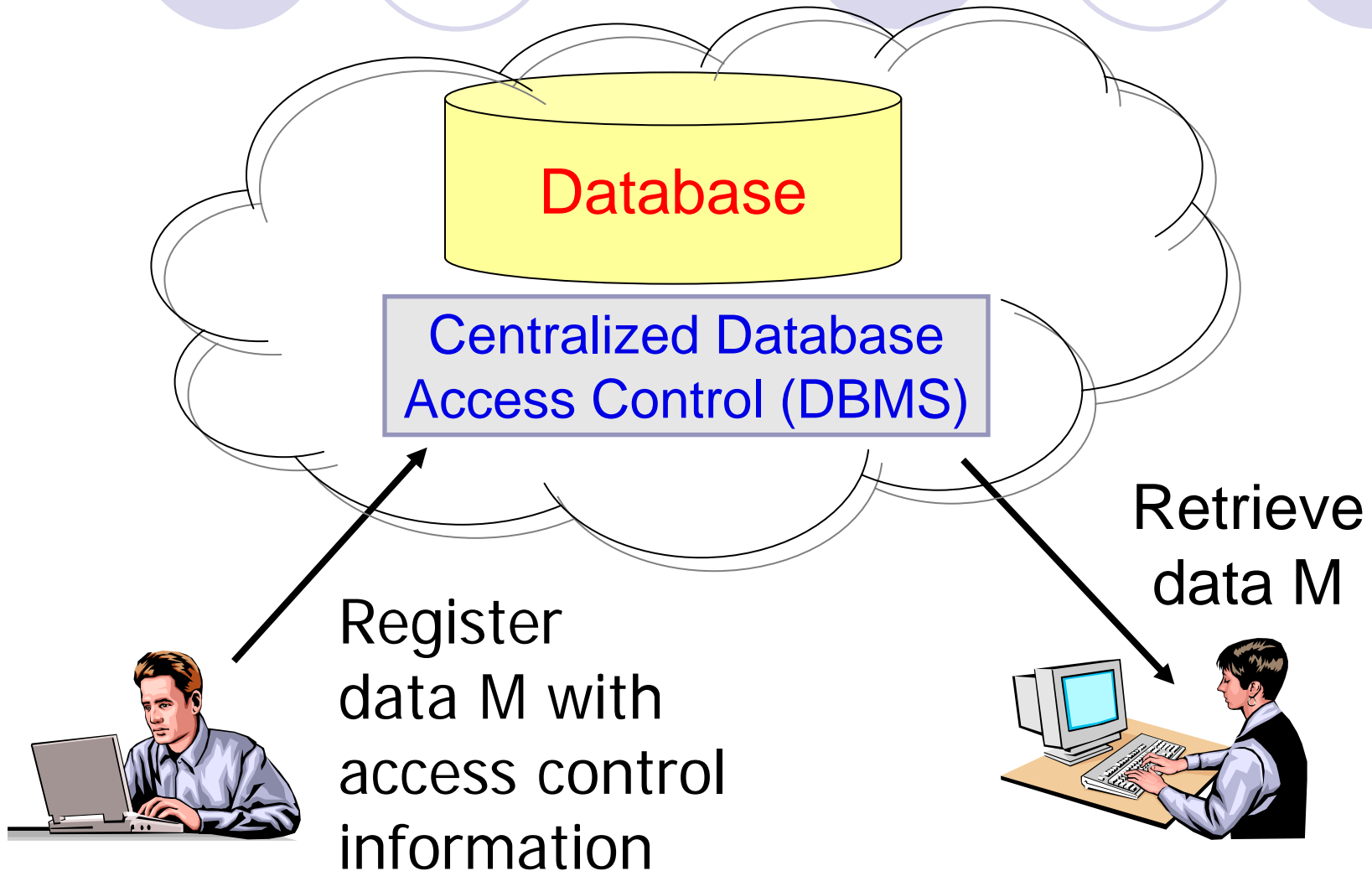
(Encryption)

C

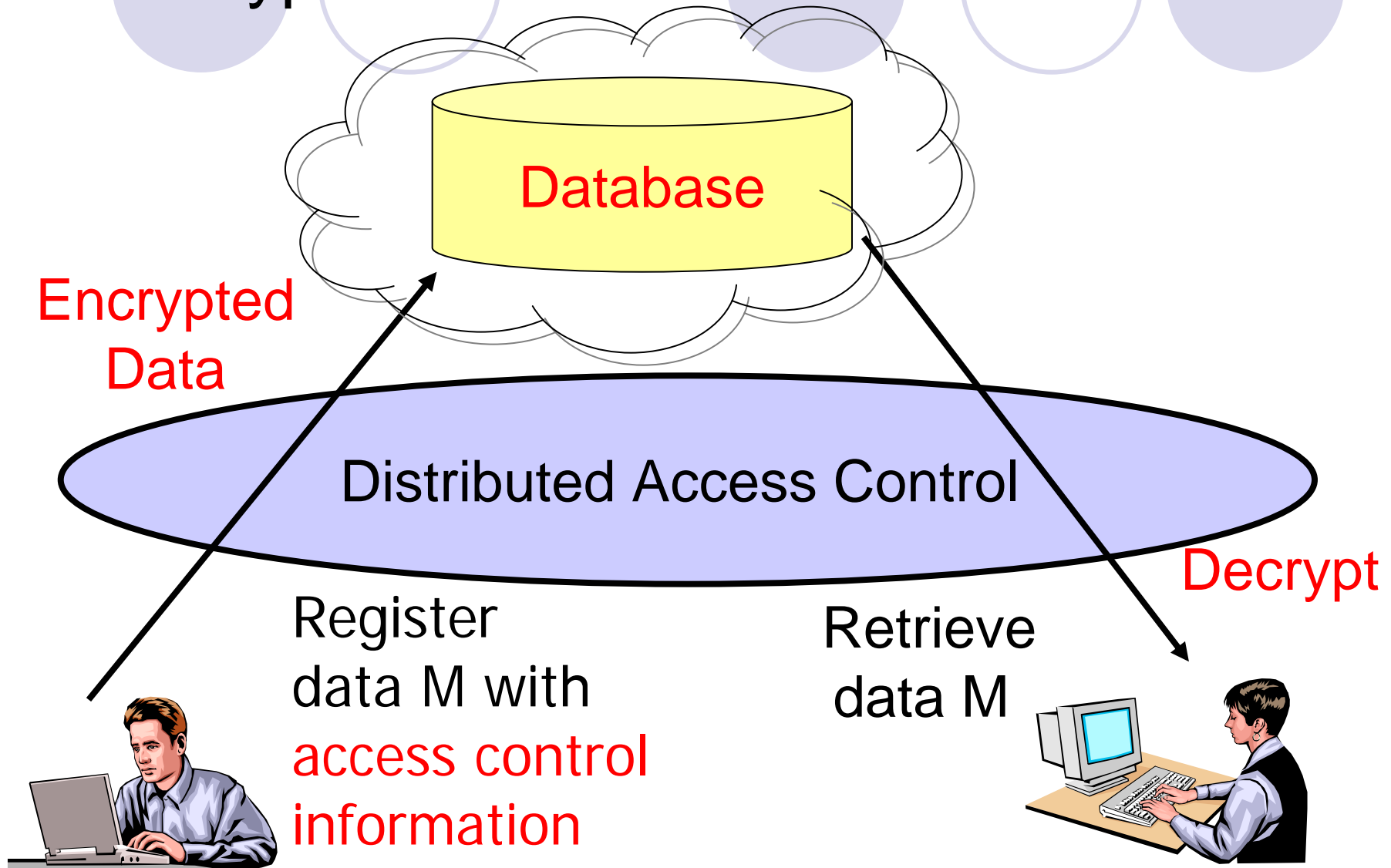
$$M = D(SK_{Alice}, C)$$

(Decryption)

Centralized Access Control of Database Services



Distributed Access Control with Encryption for Database Services



Fine-Grained Encryption (Predicate Encryption)

Public-Key
(system parameters)

PK

Publish



Authority

Public-Key: PK
Secret-Key:
 SK
(Master- SK)

Download

Bob



Predicate f

Secret-Key
for f
(SK_f)



Alice

$$C = E(x, PK, M)$$

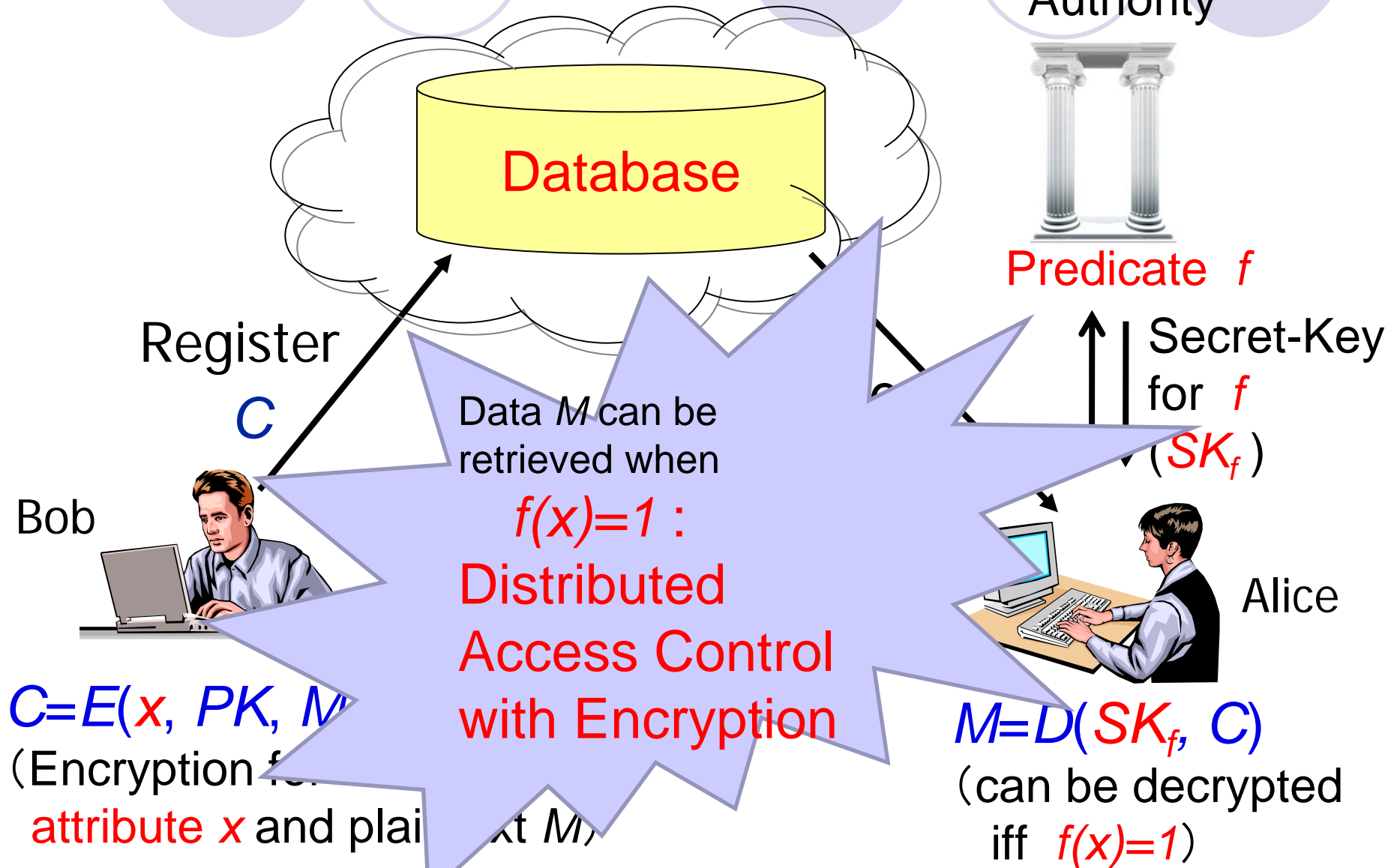
(Encryption for
attribute x and plaintext M)

C

$$M = D(SK_f, C)$$

(can be decrypted
iff $f(x) = 1$)

Distributed Access Control by Fine-Grained Encryption



Example 1 of Predicates and Attributes

- Predicate f with parameters a :

$$f(X) \equiv (X=a)$$

- Acceptable attribute x is a ,
- Application: Identity-based encryption

$a = \text{Alice@ibe.com}$

$SK_{\text{Alice}} = SK_a$



Alice

$C = E(a, PK, M)$
(Encryption)



$M = D(SK_a, C)$
(Decryption)

Example 2 of Predicates and Attributes

- Predicate f with parameters a , b and c :

$$f(X, Y, Z) \equiv (X=a) \wedge ((Y=b) \vee (Z=c))$$

- Acceptable attributes: $x = (a, b, d)$, (a, e, c)

- Application Example:

$a \equiv$ [Type: Animation],

$b \equiv$ [Price: Zone 2],

$c \equiv$ [Restriction: Class 1],

- Secret-key SK_f is given to a person (receiver) who purchases the secret key.
- A ciphertext of a content with attribute (Animation, Price Zone 2, Class 2) or (Animation, Price Zone 1, Class 1) can be decrypted by SK_f .

Fine-Grained Encryption (Predicate Encryption)

Public-Key
(system parameters)

PK

Publish



Public-Key: PK
Secret-Key:
 SK
(Master- SK)



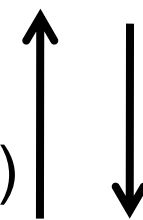
Authority

Download

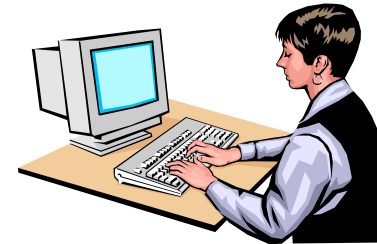
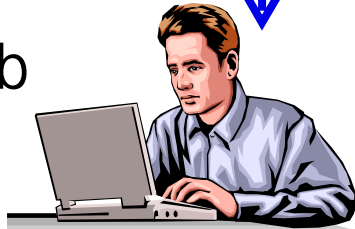


Predicate $f =$
 $(X=Animation) \wedge$
 $(Y=Price\ Zone\ 2 \vee Z=Class\ 1)$

Secret-Key
for f
 (SK_f)



Bob



Alice

$$C = E(x, PK, M)$$

(Encryption for

attribute $x = (X, Y, Z) =$

(Animation, Price Zone 2, Class 2))

C



$$M = D(SK_f, C)$$

(can be decrypted

iff $f(x)=1$)

Example 3 of Predicates and Attributes

- Predicate f with parameters a , b and c :

$$f(X, Y, Z) \equiv (X \neq a) \wedge ((Y < b) \vee (Z = c))$$

- Acceptable attributes: $x = (c, d, e)$ (s.t. $c \neq a$ and $d < b$),
 (c, e, c) (s.t. $c \neq a$ and $e > b$), ...

- Application Example:

$a \equiv$ [Type: Animation],

$b \equiv$ [Price: \$100],

$c \equiv$ [Restriction: Class 1],

- Secret-key SK_f is given to a person (receiver) who purchases the secret key.
- A ciphertext of a content with attribute (Drama, \$70, Class 2) or (News, \$200, Class 1) can be decrypted by SK_f .

Fine-Grained Encryption (Predicate Encryption)

Public-Key
(system parameters)

PK

Publish



Public-Key: PK
Secret-Key:
 SK
(Master- SK)



Authority

Download

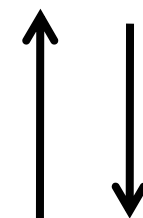


Bob



Predicate $f =$
 $(X \neq \text{Animation}) \wedge$
 $(Y < \$100 \vee Z = \text{Class } 1)$

Secret-Key
for f
(SK_f)



Alice

$$C = E(x, PK, M)$$

(Encryption for

attribute $x = (X, Y, Z) =$
(Drama, \$70, Class 2))



C

$$M = D(SK_f, C)$$

(can be decrypted
iff $f(x) = 1$)

Example 4 of Predicates and Attributes

- **Predicate f** with parameters a, b, c and d :
$$f(X, Y, Z, W) \equiv \text{TH2} [(X = a), ((Y < b) \vee (Z = c)), (W \neq d)]$$

(TH2: threshold function that accepts if at least 2 terms are true.)
- Acceptable attributes: $x = (a, e, g, d)$ (s.t. $e < b$), (d, e, c, g) (s.t. $d \neq a, e > b, g \neq d$), ...
- Application Example:
 $a \equiv [\text{Type: Animation}],$ $b \equiv [\text{Price: \$100}],$
 $c \equiv [\text{Membership: Gold}],$ $d \equiv [\text{Restriction: Class 1}]$
- Secret-key SK_f is given to a person (receiver) who purchases the secret key.
- A ciphertext of a content with attribute (Animation, \$70, Silver, Class 1) or (News, \$200, Gold, Class 2) can be decrypted by SK_f .

Fine-Grained Encryption (Predicate Encryption 2)

Public-Key
(system parameters)

PK

Publish



Authority

Public-Key: PK
Secret-Key:
 SK
(Master- SK)

Download

Bob



attribute x

Secret-Key
for x
(SK_x)



Alice

$$C = E(f, PK, M)$$

(Encryption for
predicate f and plaintext M)

C

$$M = D(SK_x, C)$$

(can be decrypted
iff $f(x) = 1$)

Fine-Grained Encryption (Predicate Encryption 2)

Public-Key
(system parameters)

PK

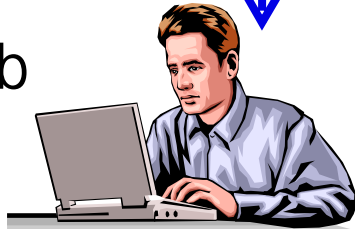
Publish



Download



Bob



$$C = E(f, PK, M)$$

(Encryption for

predicate $f =$

$$X = \text{NTT} \wedge (Y < 35 \vee Z = \text{Male})$$

C



$$M = D(SK_x, C)$$

(can be decrypted

iff $f(x) = 1$)

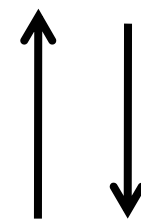
Public-Key: PK
Secret-Key:
 SK
(Master- SK)



Authority

Attribute $x = (X, Y, Z) =$
(NTT, Age:30, Female)

Secret-Key
for x
(SK_x)



Alice

From Coarse-Grained to Fine-Grained Encryption

- **Symmetric Encryption (ancient ~)**
 - sender and receiver should be fixed/restricted.
- **Public-Key Encryption (1976 ~) DH**
 - receiver should be restricted by the public-key (registered string), but sender is not restricted.
- **Identity-Based Encryption (1984/2000~) S, SK/BF, ..**
 - senders and receivers are related by an arbitrary string (e.g., identity), and sender is not restricted.
- **Predicate/Attribute-Based Encryption (2005~) SW, ..**
 - senders and receivers are only related by logical conditions (predicate-attribute relations), and sender is not restricted.

Mathematics to Realize Cryptosystems

Permutations
, Substitutions
(Ancient~)

Cyclic Groups
(1976~)

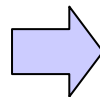
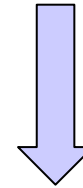
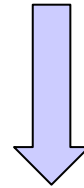
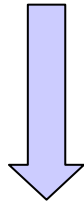
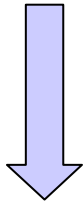
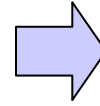
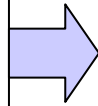
Bilinear Pairing
Groups
(2000~)

Symmetric
Encryption
(Ancient~)

Public-key
Encryption
(1976~)

ID-Based
Encryption
(2000~)

Attribute-Based/
Predicate
Encryption
(2005~)



Cryptography and the Internet

Public-key
cryptosystem

RSA scheme

Attribute-Based /
Predicate
Encryption

Identity-Based
Encryption

IPSEC, SSL
PKI

1969

1976

1983

1988

1991

1993

2000

Present

Arpanet (the origin
of the Internet) started

Commercial usage
started

World Wide Web
by CERN

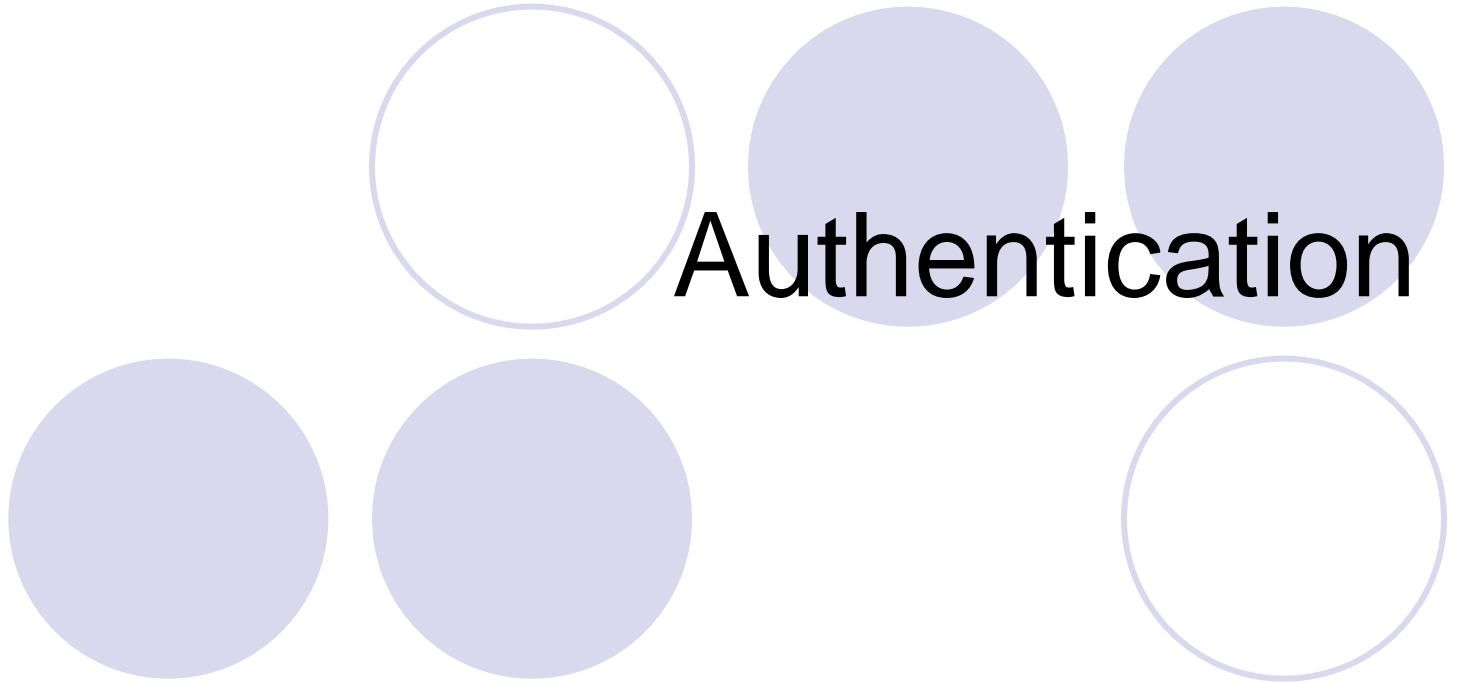
Browser (Mosaic)

TCP/IP protocol
(changed from original NCP)

Web 2.0

Cloud
computing

Authentication

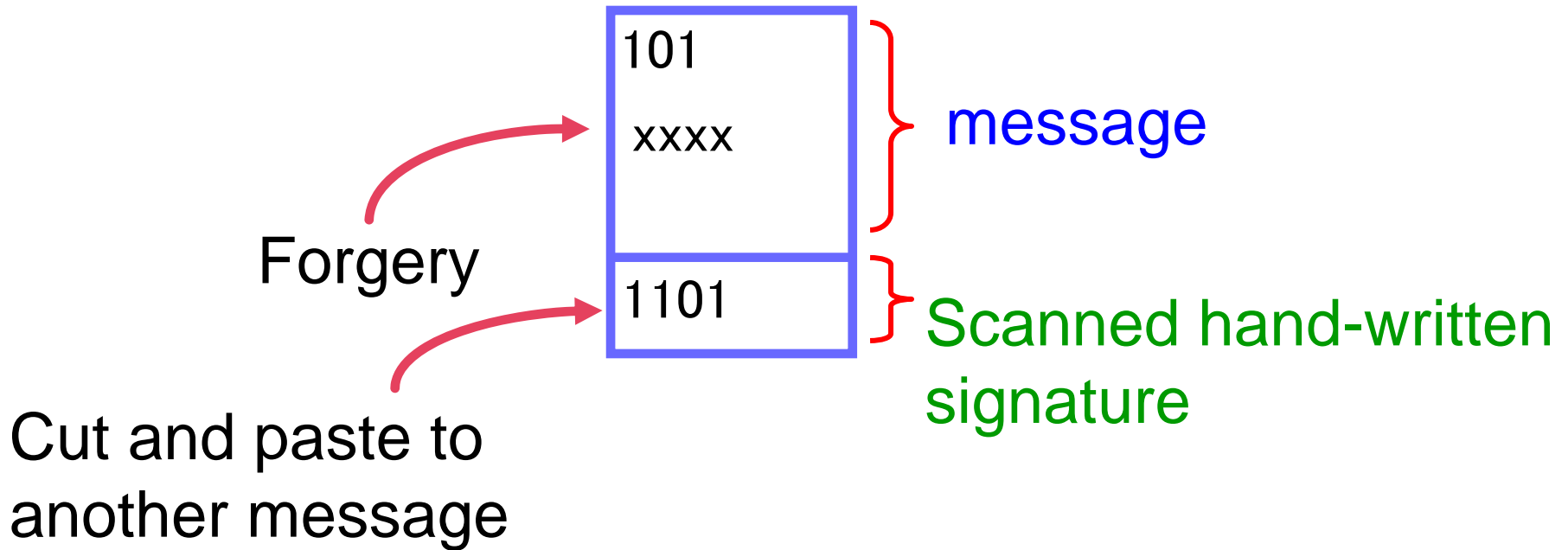


Another problem raised in the early 1970's

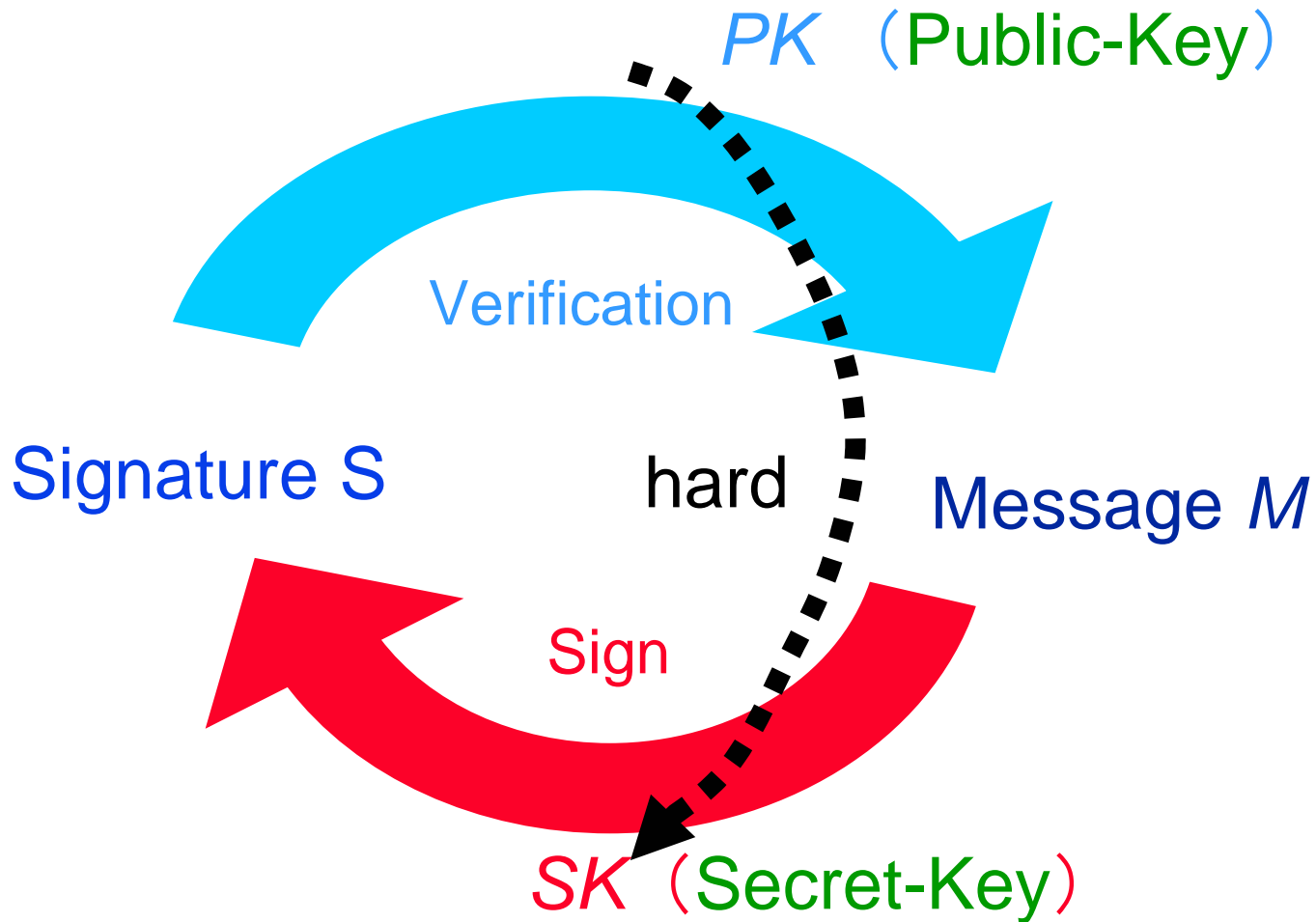
- **Problem:** Is it possible to sign on an electronic (digital) message?

Difficulty of Realizing Electronic Signatures

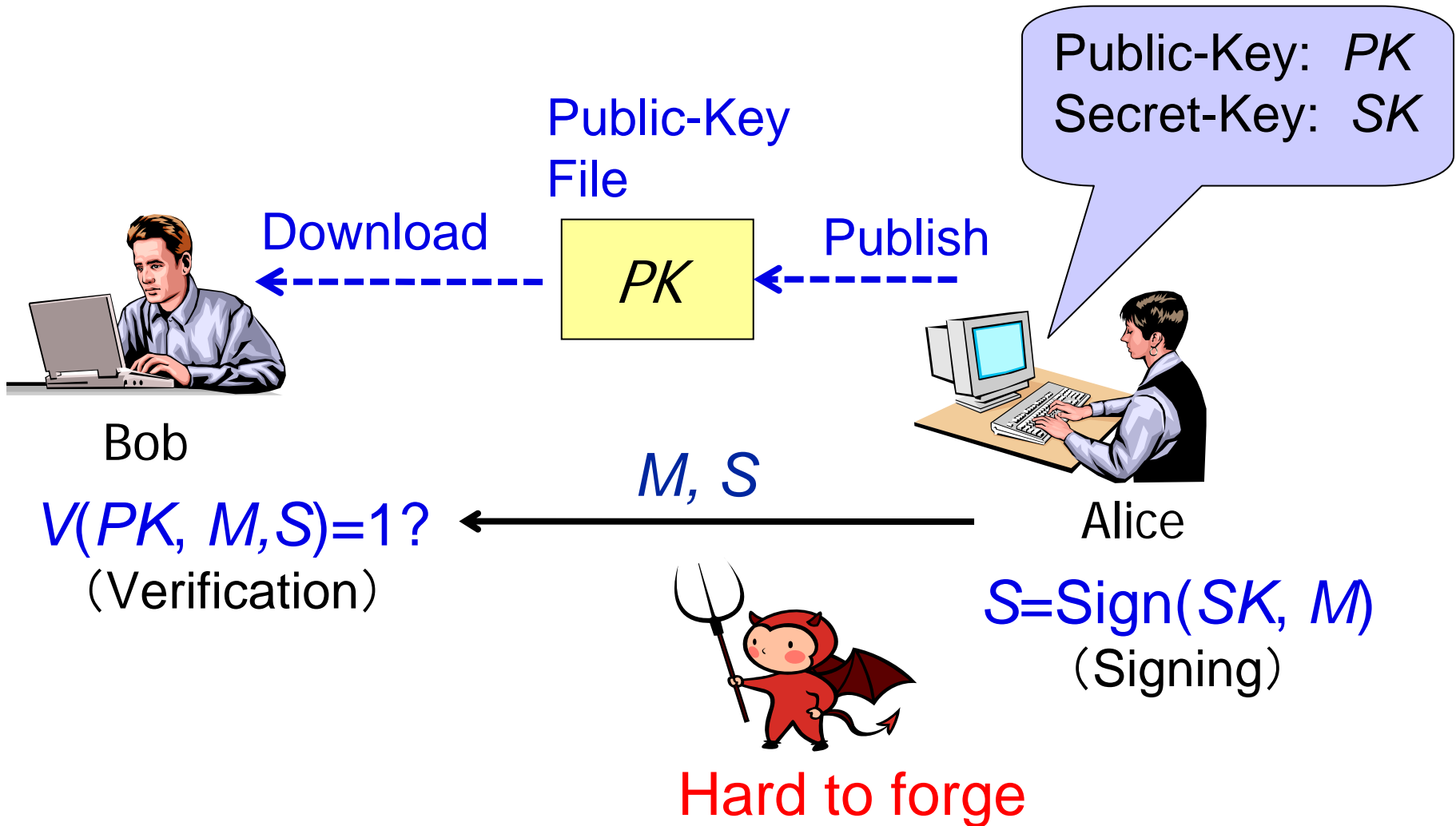
Easy to forge a digital message with no mark



Principle of Digital Signatures



Digital Signatures



Public-Key Infrastructures (PKI) Certification Authority (CA)

Organization to certify a public-key

Secret-Key Public-Key



Register Public-Key PK

Certification
Authority (CA)



User A



Certifying PK
with User A



CA's Sign

After confirming the
identity of User A,
issues the certification

The slide features five circles of varying shades of light purple. Two are solid, and three are hollow. The text is centered over the hollow circles in the upper right quadrant.

Privacy-Enhanced Authentication



Privacy-Enhanced Signatures

- **PKI-based digital signatures:**

The signer's identity of signatures is rigorously confirmed. It is good for the standard usage of applications of signatures.

- In some applications, it is not good, where **the signer should be anonymous, and some qualification (e.g., a member of a group) of the signer should be authenticated.**

Group Signatures

Group Public-Key

PK

Publish

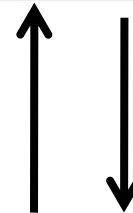


Public-Key: PK
Secret-Key:
 SK
(Master- SK)

Authority
(Group Manager)

Secret-Key
for *membership*
(SK_{Group})

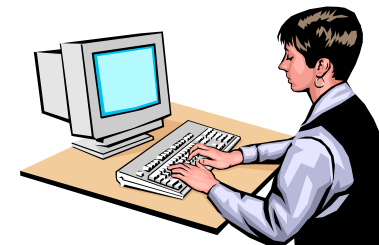
Membership



Download



Bob



Alice

$V(PK, M, S) = 1?$
(Verification)

M, S

$S = \text{Sign}(SK_{Group}, M)$
(Signing)

Fine-Grained (Predicate) Signatures

Public-Key
(System parameters)

PK

Publish

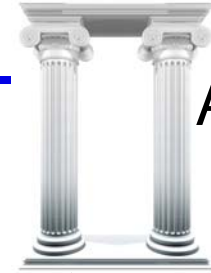
Download

Bob



$V(PK, M, S, f) = 1?$
(Verification)

Public-Key: PK
Secret-Key:
 SK
(Master- SK)



Authority

Attribute x

Secret-Key
for *attribute* x
(SK_x)



Alice

$S = \text{Sign}(SK_x, f, M)$
(Signing for f
s.t. $f(x) = 1$)

Example of Predicates and Attributes

- **Attribute** $x = (X, Y, Z)$
= ([Group: **Group 2**], [Age: **30**], [Gender: **Female**])
- Secret-key SK_x is given to a person who has the attributes.
- **Predicate** f with parameters a , b and c :
 $f(X, Y, Z) \equiv (X = a) \wedge ((Y < b) \vee (Z = c))$, where
 $a \equiv$ [Group: **Group 2**],
 $b \equiv$ [Age: **35**],
 $c \equiv$ [Gender: **Male**]
- A signature with predicate f implies that the attributes of the signer satisfies the predicate f ,
i.e., $(X = \text{Group 2}) \wedge ((Y < 35) \vee (Z = \text{Male})) = 1$ (True)

Fine-Grained (Predicate) Signatures

Public-Key
(System parameters)

PK

Publish

Public-Key: PK
Secret-Key:
 SK
(Master-SK)



Authority

Download

Attribute $x = (X, Y, Z)$
(Group2, Age:30, Female)

Secret-Key
for *attribute* x
(SK_x)

Bob



Alice

$V_{PK}(M, S, f) = 1?$
(Verification)

$M, S, \text{predicate } f =$

$(X = \text{Group 2}) \wedge$
 $((Y < \text{Age:35}) \vee (Z = \text{Male}))$ s.t. $f(x) = 1$

$S = \text{Sign}(SK_x, f, M)$
(Signing for f)

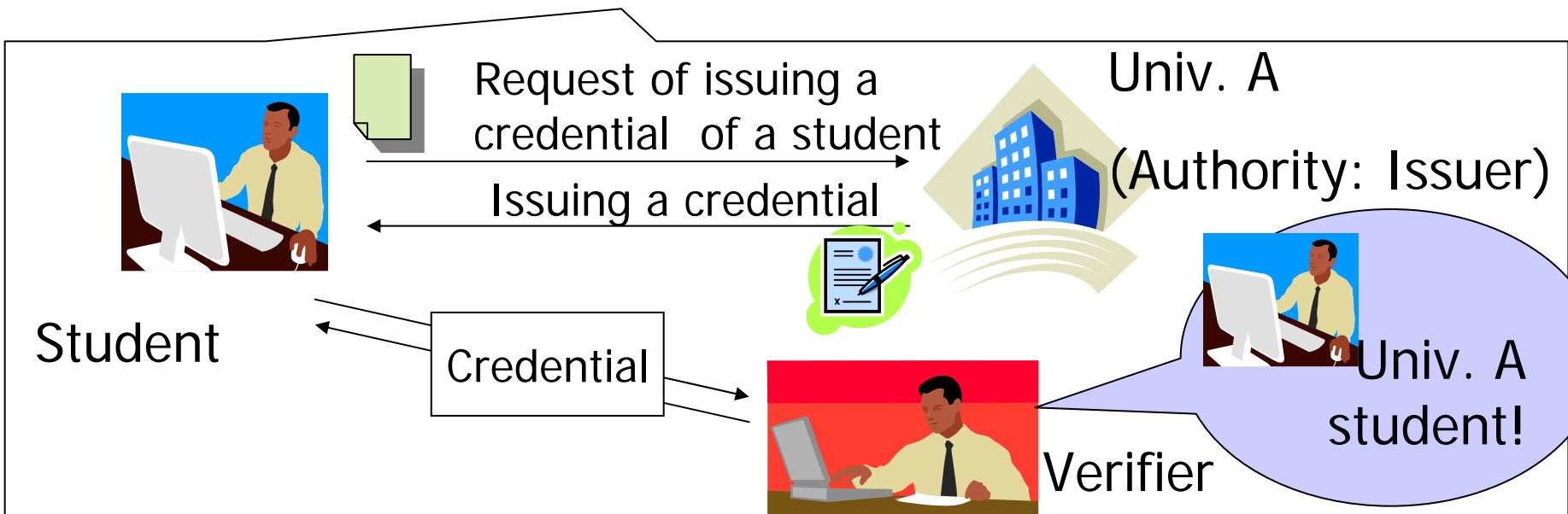
Example of Predicates and Attributes

- **Attribute** $x = (X, Y, Z, W) =$
([Group: **Group 1**], [Age: **20**], [Gender: **Male**], [Member: **Silver**])
- Secret-key SK_x is given to a person who has the attributes.
- **Predicate** f with parameters a, b, c and d :
 $f(X, Y, Z, W) \equiv \text{TH2} [(X = a), ((Y < b) \vee (Z = c)), (W \neq d)]$
(TH2: threshold function that accepts if at least 2 terms are true.)
- Application Example:
 $a \equiv$ [Group: **Group 2**], $b \equiv$ [Age: **30**],
 $c \equiv$ [Gender: **Female**], $d \equiv$ [Member: **Gold**],
- A signature with predicate f implies that the attributes of the signer satisfies the predicate f , i.e.,
 $\text{TH2} [(X = \text{Group 2}), ((Y < 30) \vee (Z = \text{Female}))], (W \neq \text{Gold})] = 1.$

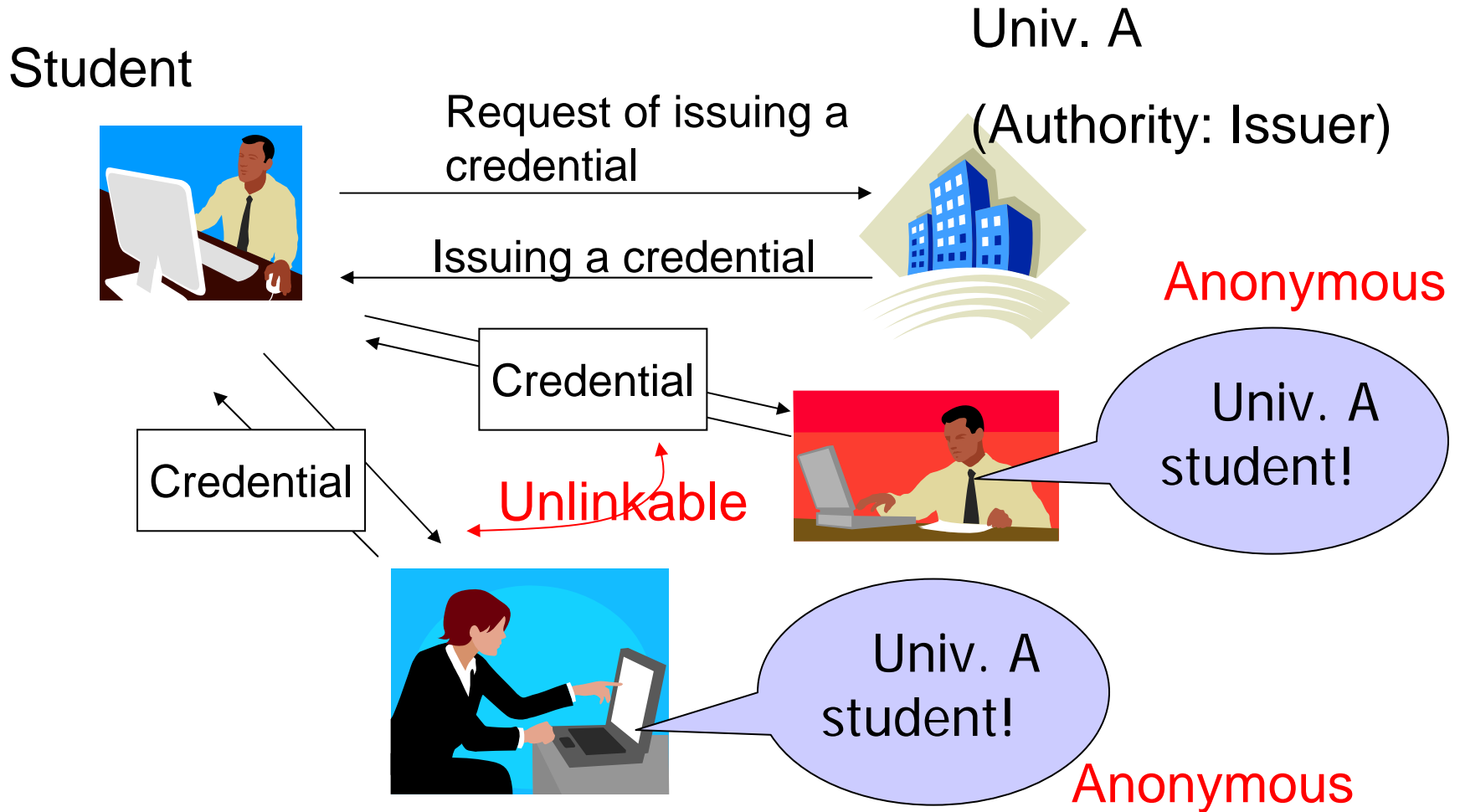
Credential

Certificate for person's qualification/attribut

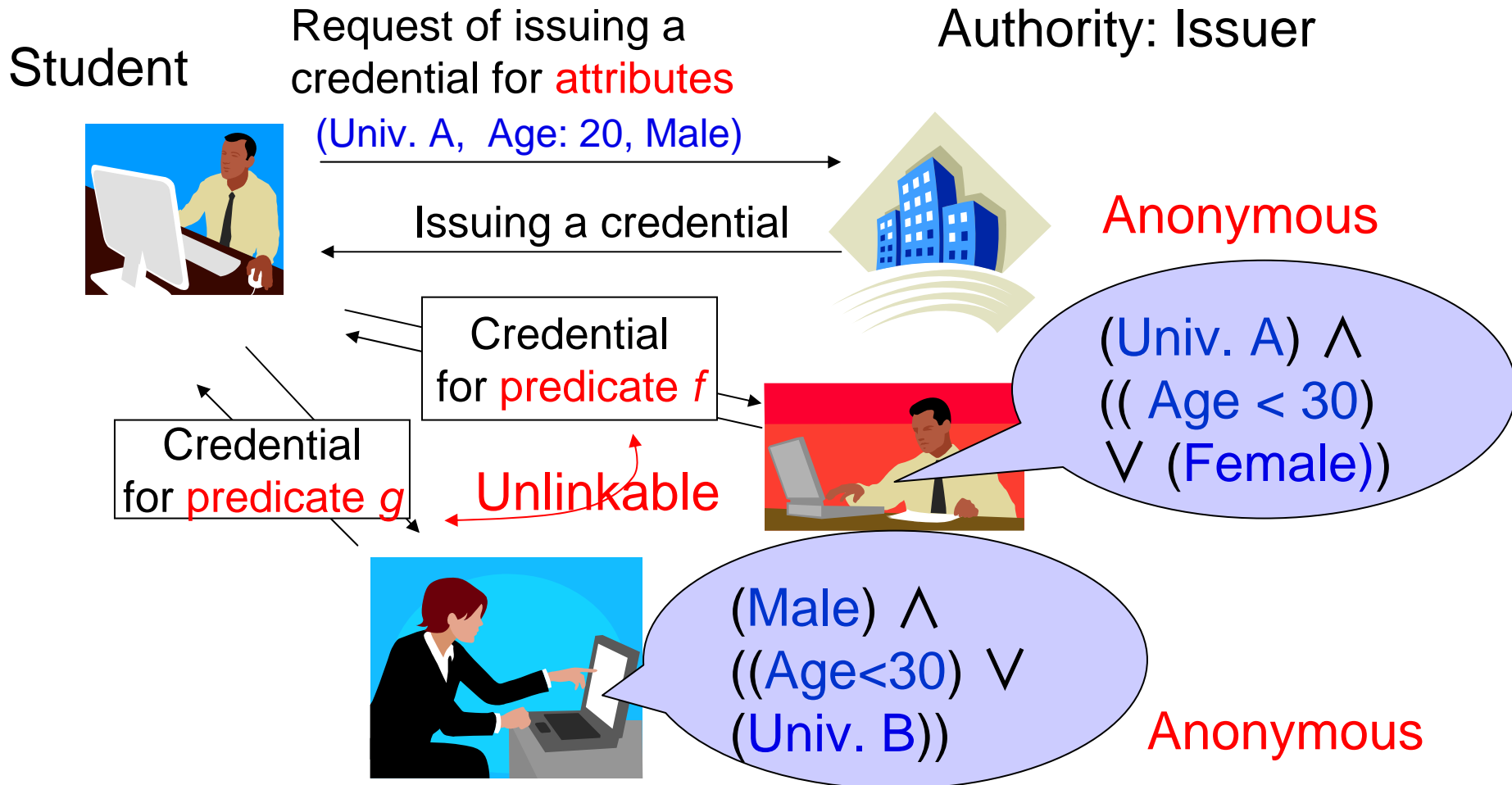
Eg.) "Student of University A", "Right to enter a room"



Anonymous Credential



Fine-Grained Anonymous Credential



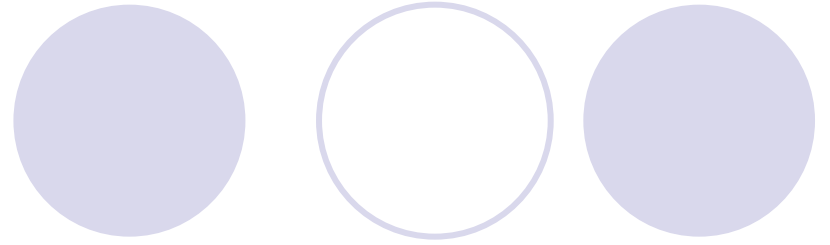
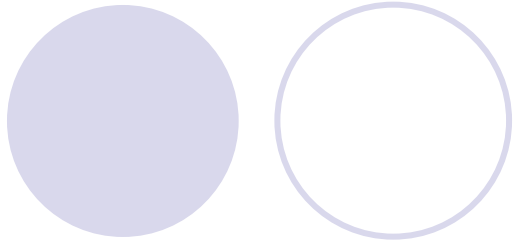
For Applications of Fine-Grained Crypto to Practice

- Some standard frameworks (classifications and formats) defining attributes and predicates
- Management and usage rules and tools:
New infrastructures (like PKI)
- It took around 20 years to establish PKI after the concept of public-key cryptosystem was proposed. It will take a certain period for the fine-grained cryptosystems to be widely used in practice.

For Your Further Study

Five circles are arranged horizontally at the top of the slide. From left to right: a solid light purple circle, an outlined light purple circle, a solid light purple circle, an outlined light purple circle, and a solid light purple circle.

- IACR (International Association for Cryptologic Research) ePrint Archive
<http://www.iacr.org/>
- Key Words:
 - Predicate Encryption
 - Attribute-Based Encryption (ABE)
 - Attribute-Based Signatures (ABS)



Thank you !!