

Linear-Time and Efficient Distributed Algorithms for List Coloring Graphs on Surfaces

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Abstract—In 1994, Thomassen proved that every planar graph is 5-list-colorable. In 1995, Thomassen proved that every planar graph of girth at least five is 3-list-colorable. His proofs naturally lead to quadratic-time algorithms to find such colorings. Here, we provide the first linear-time algorithms to find such colorings.

For a fixed surface S , Thomassen showed in 1997 that there exists a linear-time algorithm to decide if a graph embedded in S is 5-colorable and similarly in 2003 if a graph of girth at least five embedded in S is 3-colorable. Using the theory of hyperbolic families, the author and Thomas showed such algorithms exist for list-colorings. Around the same time, Dvorak and Kawarabayashi also provided such algorithms. Moreover, they gave an algorithm to find such colorings (if they exist). Here we provide the first such algorithm which is fixed parameter tractable with genus as the parameter; indeed, we provide a linear-time algorithm to find such colorings.

In 1988, Goldberg, Plotkin and Shannon provided a deterministic distributed algorithm for 7-coloring n -vertex planar graphs in $O(\log n)$ rounds. In 2018, Aboulker, Bonamy, Bousquet, and Esperet provided a deterministic distributed algorithm for 6-coloring n -vertex planar graphs in polylogarithmic rounds. Their algorithm in fact works for 6-list-coloring. They also provided a polylogarithmic algorithm for 4-list-coloring triangle-free planar graphs. Chechik and Mukhtar independently obtained such algorithms for ordinary coloring in $O(\log n)$ rounds, which is best possible in terms of running time. Here we provide the first polylogarithmic deterministic distributed algorithms for 5-coloring n -vertex planar graphs and similarly for 3-coloring planar graphs of girth at least five. Indeed, these algorithms run in $O(\log n)$ rounds, work also for list-colorings, and even work on a fixed surface (assuming such a coloring exists).

Index Terms—Graph Coloring; Distributed Algorithms; Planar Graphs; List Coloring

I. INTRODUCTION

List Coloring Planar Graphs. Graph coloring is a widely studied area in graph theory. In particular, coloring planar graphs - and more generally coloring graphs embedded in surfaces - has received much attention. Recall that a k -coloring of a graph G is an assignment of colors $\{1, \dots, k\}$ such that adjacent vertices do not receive the same color. In 1977, Appel and Haken [1], [2] proved the Four Color Theorem that every planar graph has a 4-coloring. In 1989, they [3] gave a quartic-time algorithm to find such a coloring, which was improved to a quadratic-time algorithm by Robertson, Sanders, Seymour and Thomas [24] in 1996. In 1959, Grötzsch proved that every triangle-free planar graph G has a 3-coloring. Thomassen's short proof [29] of Grötzsch's theorem can easily be converted

to a quadratic-time algorithm to find such a coloring. This was improved to $O(v(G) \log v(G))$ by Kowalik [13] in 2004 and finally to a linear-time algorithm in 2009 by Dvořák, Kawarabayashi and Thomas [9].

List-coloring is a generalization of coloring introduced by Erdős, Rubin and Taylor [10] in 1979 and independently by Vizing [30] in 1976. A *list-assignment* L of a graph G is an assignment of lists of colors $L(v)$ to each vertex v of G ; we say L is a k -list-assignment if $|L(v)| \geq k$ for every vertex v . An L -coloring ϕ is a coloring such that $\phi(v) \in L(v)$ for each vertex v . We say a graph G is k -list-colorable if G has an L -coloring for every k -list-assignment L . In 1994, Thomassen [25] proved that every planar graph is 5-list-colorable; this is best possible in that Voigt [31] constructed a planar graph that is not 4-list-colorable. In 1995, Thomassen [26] proved that every planar graph having girth at least five is 3-list-colorable; this is best possible as Voigt [32] constructed a triangle-free planar graph that is not 3-list-colorable. Of course, every triangle-free planar graph is 4-list-colorable as such graphs are 3-degenerate.

For the sake of convenience when stating theorems, we say a list assignment L of a graph G is *type 345* if L is a k -list-assignment for some $k \in \{3, 4, 5\}$ and G has girth at least $8 - k$. Thomassen's proofs naturally lead to quadratic-time algorithms to find L -colorings for type 345 list-assignments. It seems conceivable that a clever algorithmic implementation of Thomassen's proofs would yield linear-time algorithms; however, no such algorithm has appeared in the literature to date.

Our first main result (a special case of our more general result about surfaces stated later) is to provide the first linear-time algorithms to find such colorings.

Theorem 1.1. *There exists a linear-time algorithm to find an L -coloring of a planar graph G if L is a type 345 list assignment for G .*

Coloring and List-Coloring Graphs on Surfaces. It is natural to wonder if the results mentioned above can be extended to graphs embedded in a fixed surface. Of course, there do exist graphs that are not 5-list-colorable. Since coloring is a monotone property, it is natural to study minimally non-colorable graphs. To that end, we say a graph G is *critical for k -coloring* if G is not k -colorable but every proper subgraph of

G is. Similarly we say a graph G is *critical for k -list-coloring* if there exists a k -list-assignment L such that G does not have an L -coloring but every proper subgraph of G does.

In a lengthy breakthrough work in 1997, Thomassen [27] proved that for every surface Σ , there exist only finitely many graphs embeddable in Σ that are critical for 5-coloring. As a corollary, it follows that there exists a linear-time algorithm for each surface Σ to decide if a graph embeddable in Σ is 5-colorable. Similarly in 2003 in yet another deep work, Thomassen [28] proved that for every surface Σ , there exist only finitely many graphs of girth at least five embeddable in Σ that are critical for 3-coloring. Hence there exists a linear-time algorithm for each surface Σ to decide if a graph of girth at least five embeddable in Σ is 3-colorable.

Thomassen [27] naturally asked if such results could be generalized to list-coloring. For 5-list coloring, the author and Thomas [23] proved the analogous result while the author [19] proved the analogous result for 3-list-coloring graphs having girth at least five. These proofs are rather lengthy and rely on the theory of strongly hyperbolic families developed by the author and Thomas in [23]. Moreover, we proved that such critical graphs have at most $O(g(\Sigma))$ vertices, where $g(\Sigma)$ denotes the *Euler genus* of Σ (that is $2h + c$ where h is the number of handles of Σ and c is the number of crosscaps). These results imply linear-time algorithms to decide if graphs embeddable in a fixed surface are 5-list-colorable or 3-list-colorable if they have girth at least five. Indeed, our algorithms can also decide if an L -coloring exists for a fixed type 345-list-assignment L .

As for finding such an L -coloring (if it exists) given a fixed type 345 list-assignment L , Dvořák and Kawarabayashi [8] used some of the results on hyperbolic families combined with some results about the treewidth of graphs on surfaces to provide an $O(v(G)^{O(g(\Sigma)+1)})$ -time algorithm to find an L -coloring of a graph embedded in Σ , provided such a coloring exists. Dvořák and Kawarabayashi [8] asked whether the dependence on $g(\Sigma)$ in the exponent could be removed. This would imply that finding a 5-list-coloring of a graph (if it exists) and similarly finding a 3-list-coloring for a girth five graph (if it exists) is fixed parameter tractable (FPT) where the parameter is the genus. We answer this question in the affirmative. Indeed, we generalize Theorem I.1 as follows.

Theorem I.2. *For each surface Σ , there exists a linear-time algorithm to find an L -coloring (if it exists) of a graph G embeddable in Σ if L is a type 345 list assignment for G .*

We note that there are many graphs on surfaces which do have an L -coloring for every type 345 list assignment. In particular, the author and Thomas [23] showed that graphs with edge-width $\Omega(\log g(\Sigma))$ have this property.

Deterministic Distributed Algorithms for Coloring and List-Coloring Planar Graphs. Another paradigm of studying for coloring algorithms is distributed algorithms. Here we study the synchronous message-passing model of distributed computing, in particular the LOCAL model (where unlim-

ited size messages are allowed). However, we believe our algorithm should also work in the CONGEST model (where there are limits on message sizes) but do not pursue this direction of study. Moreover, we restrict our attention to the study of deterministic algorithms in the LOCAL model. In these models, round complexity is the benchmark. Indeed, efficiency is usually defined as having round complexity at most polylogarithmic in the number of vertices. We refer the reader to the survey book of Barenboim and Elkin [6] for more details on distributed computing.

What then is known for deterministic distributed algorithm for coloring planar graphs? In 1988, Goldberg, Plotkin and Shannon [12] provided a deterministic distributed algorithm for 7-coloring n -vertex planar graphs in $O(\log n)$ rounds. It should be noted that Goldberg, Plotkin and Shannon gave an $O(\log n \log^* n)$ round parallel algorithm for 5-coloring n -vertex planar graphs. They falsely claimed that all of the results in that paper carried over to the distributed setting but this turns out not to be the case for 5-coloring planar graphs. Rather their algorithm for 5-coloring planar graphs would give $O(n)$ round complexity in the distributed setting. This inaccuracy seems to have been perpetuated (e.g. in [5]) until the authors in [4] noticed the flaw. Hence the question of improving the distributed result for 5-coloring and even 6-coloring planar graphs had in fact been open.

In 2018, Aboulker, Bonamy, Bousquet, and Esperet [4] provided a deterministic distributed algorithm for 6-coloring n -vertex planar graphs in $O(\log^3 n)$ rounds. Their algorithm in fact works for 6-list-coloring. They also provided an $O(\log^3 n)$ -round algorithm for 4-list-coloring triangle-free planar graphs. Chechik and Mukhtar [7] independently obtained such algorithms for ordinary coloring in $O(\log n)$ rounds, which is best possible in terms of round complexity. Aboulker et al. [4] asked whether such efficient algorithms exist for 5-coloring planar graphs and 3-coloring planar graphs of girth at least five.

Aboulker et al. [4] proved that no such efficient algorithms exist for 4-coloring planar graphs and 3-coloring triangle-free planar graphs. This is because such algorithms would imply that graphs whose $\Omega(\log n)$ neighborhoods are planar are 4-colorable (and 3-colorable if triangle-free) and there exist graphs which do not satisfy these properties. They thus also noted that their conjectures for the remaining cases would then imply that $\Omega(\log n)$ locally planar graphs are 5-colorable and 3-colorable if girth at least five. That result is true but the only known proof of it follows from the author and Thomas' general result [23] about hyperbolic families.

Here we provide the first polylogarithmic deterministic distributed algorithms for 5-coloring n -vertex planar graphs and similarly for 3-coloring n -vertex planar graphs of girth at least five. Indeed, these algorithms run in $O(\log n)$ rounds and work even for list-colorings. In addition, they even work on a fixed surface assuming a coloring exists.

Theorem I.3. *For each surface Σ , there exists a deterministic distributed algorithm that given a graph G embeddable in Σ*

and a type 345 list-assignment L for G , finds an L -coloring of G (if it exists) in $O(\log v(G))$ rounds.

II. PROOF OVERVIEW AND ALGORITHM

In this section, we first provide an overview of the proofs of Theorems I.2 and I.3 at a high level. Then in the remaining subsections of this section, we proceed to give the formal definitions and statements of the major theorems used in their proofs. We also state the algorithm used for both Theorem I.2 and Theorem I.3. Finally, we provide an outline of the rest of the paper at the end of the section.

Proof Overview. We provide the same algorithm (Algorithm 1, stated in Section II-C) for both Theorems I.2 and I.3. The complexity analysis for the two theorems is of course different but in both cases is rather straightforward (see Section III for details). The proof of the correctness, however, relies on the deep theorems proved by the author and Thomas described in the introduction. Nevertheless, there are two key new contributions in this work.

The first contribution is a structure theorem (Theorem II.3) for planar graphs. Theorem II.3 roughly says that a planar graph either has a 2-separation with a constant-size side, or has a constant-size subgraph that is ‘deep’ (i.e. the ratio between its vertices and its neighbors is large). More than that, the theorem guarantees linearly many pairwise non-touching such subgraphs. We note that Theorem II.3 is not about coloring and hence may be of independent interest for developing efficient algorithms. Indeed, this structure theorem is essential for establishing the running time in both algorithms.

The second contribution is to combine this structure theorem with the deep coloring theorems proved elsewhere to prove a startling new coloring theorem (Theorem II.7). Theorem II.7 roughly says that for every planar graph G , there exists a constant-size, constant-degree subgraph that is ‘deletable’ for type 345 list assignments (i.e. every coloring of the remaining graph extends to a coloring of the subgraph no matter the list assignment). More than that, the theorem guarantee linearly many pairwise non-touching such subgraphs. We note that our notion of deletable (see Definition II.5) is quite strong, essentially the strongest such notion one could define.

The algorithm then consists in finding and deleting such deletable subgraphs, then recursively finding a coloring of the smaller graph and finally extending that coloring to the deleted subgraphs. Since these subgraphs are constant-size and constant-degree these searches, deletions and extensions can be performed very efficiently. Since there are linearly many such subgraphs, the smaller graph will be proportionally smaller (i.e. $(1 - \varepsilon)v(G)$ for some fixed ε). Combined these two facts lead to a linear-time algorithm (or a $O(\log n)$ -round algorithm in the distributed case).

We note that in the distributed case some care is needed when extending the colorings to guarantee $O(\log n)$ -round complexity as opposed to $O(\log n \log^* n)$ -round complexity. Namely that while the algorithm recurses, the deleted vertices must compute a constant coloring of an auxiliary

graph on the deletable subgraphs needed for extending the coloring efficiently. Such a coloring can be computed in $\tilde{O}(\sqrt{\Delta(G)}) + O(\log^* n)$ -round (where $\Delta(G)$ denotes the maximum degree of G) using the algorithm of Fragniaud, Heinrich and Kosowski [11]. Since $\Delta(G)$ will be a constant in our application, this will run in $O(\log^* n)$ time. To avoid $O(\log n \log^* n)$ -round complexity, though, it is crucial that this coloring be computed while the algorithm recurses so that the overall complexity is $O(\log n) + O(\log^* n) = O(\log n)$ -rounds.

In the remainder of this section, we present the formal versions of Theorems II.3 and II.7 and then state Algorithm 1. In the remaining sections of the paper, we then present the proof of the theorems and algorithm complexity.

A. A Structure Theorem for Planar Graphs

Here we state our structure theorem for planar graphs but first some definitions.

Definition II.1. Let G be a graph, H be an induced subgraph of G and $C, k > 0$. The *boundary* of H in G is $\{v \in V(H) : N(v) \setminus V(H) \neq \emptyset\}$. The *coboundary* of H in G is $\bigcup_{v \in H} N(v) \setminus V(H)$ (i.e. the boundary of $G[V(G) \setminus V(H)]$). We say H is

- a *purse* if H is connected, the size of the coboundary of H (call it S) is at most two, and $G[S \cup V(H)] + \{uv : u \neq v \in S\}$ is planar,
- a *C -pocket* if H is connected, $v(H) \leq C$ and $d_G(v) \leq C$ for every $v \in V(H)$,
- a *C -purse* if H is a purse and a C -pocket,
- *k -deep* if the coboundary of H is non-empty and has size at most $\frac{1}{k}v(H)$.

We say two subgraphs H_1, H_2 of G are *non-touching* if $V(H_1) \cap V(H_2) = \emptyset$ and there does not exist $a_1 \in V(H_1), a_2 \in V(H_2)$ such that $a_1 a_2 \in E(G)$.

As described in the proof overview, we will prove a structure theorem that every planar graph has a C -purse or a k -deep C -pocket. More than that, we will prove it has linearly many pairwise non-touching such subgraphs. Hence the following definition.

Definition II.2. Let G be a graph and $C > 0$. A *C -wallet* of G is a set \mathcal{H} such that all of the following hold:

- every $H \in \mathcal{H}$ is a C -pocket, and
- every pair $H_1 \neq H_2 \in \mathcal{H}$ are non-touching, and
- $|\{H \in \mathcal{H}\}| \geq \frac{1}{2C}v(G)$.

Let $k > 0$. We say a C -wallet \mathcal{H} is *k -deep* if for every $H \in \mathcal{H}$, H is a purse or k -deep.

We are now prepared to state our structure theorem.

Theorem II.3. *For every $k \geq 1$, there exists $C > 0$ such that the following holds: if G is a planar graph, then there exists a k -deep C -wallet \mathcal{H} of G .*

We can extend Theorem II.3 more generally to surfaces under the assumption the graph is large as follows.

Theorem II.4. For every $k \geq 1$ and surface Σ , there exists $C > 0$ such that the following holds: if G is a graph embeddable in Σ with $v(G) \geq Cg(\Sigma)$, then there exists a k -deep C -wallet \mathcal{H} of G .

B. A Coloring Theorem

Here we state our new coloring theorem but first we define our notion of deletable as follows.

Definition II.5. We say an induced subgraph H of a graph G is r -deletable in G if for every list-assignment L of H such that $|L(v)| \geq r - (d_G(v) - d_H(v))$ for each $v \in V(H)$, H has an L -coloring.

Definition II.6. We say a C -wallet \mathcal{H} of a graph G is r -deletable if every $H \in \mathcal{H}$ is r -deletable in G .

We will prove that in graphs embedded on surfaces having girth at least $g \in \{3, 4, 5\}$, purses and k -deep pockets (for some large enough k) contain $(8 - g)$ -deletable subgraphs (Lemmas V.1 and V.2 respectively). Combined with Theorem II.4, we obtain the following result which we now state as our main coloring theorem as follows.

Theorem II.7. For each surface Σ , there exists $C > 0$ such that the following holds: If G is a graph embeddable in Σ having girth at least $g \in \{3, 4, 5\}$ and $v(G) \geq Cg(\Sigma)$, then there exists an $(8 - g)$ -deletable C -wallet of G .

Theorem II.7 is quite surprising. It implies that a minimum counterexample to the strong hyperbolicity of 5-list-coloring or 3-list-coloring graphs of girth five (e.g. for Theorem V.15) has constant size and hence that theorem can be proved by brute-force methods. It also implies that minimum counterexamples to Thomassen's theorems that planar graphs are 5-list-colorable and that planar graphs of girth at least five are 3-list-colorable have constant size and hence that those theorems can also be proved by brute-force methods (instead of clever inductive arguments).

C. A Coloring Algorithm

Before stating our algorithm, we need one definition as follows.

Definition II.8. If H is a graph and $k \geq 1$ is an integer, then H^k denotes the graph with $V(H^k) = V(H)$ and $E(H^k) = \{u \neq v \in V(H) : u \text{ is at distance at most } k \text{ from } v\}$.

We are now ready to state our algorithm for both Theorem I.2 and I.3:

Let Σ be a fixed surface, let C_0 be as in Theorem II.7 for Σ and let $C = \max\{C_0, 2\}$.

Algorithm 1: Finding an L -coloring of a graph G embedded in Σ for a type 345 list-assignment L

- 1: **if** $v(G) \leq Cg(\Sigma)$ **then** find and **return** an L -coloring of G by exhaustive search
- 2: **for** each vertex v of G of degree at most C :
Find (if one exists) an $(8 - g)$ -deletable C -pocket H_v containing v
Let $H = \bigcup_v H_v$
- 3: **recurse** on $G_0 = G - V(H)$ to find an L -coloring ψ_0 of G_0
- 4: While recursing, also find (simultaneously for LOCAL) a C^{2C} -coloring ϕ of H^{2C}
- 5: **for** each $i \in \{1, \dots, C^{2C}\}$:
(i) Let $H_i = (\bigcup_v H_v : \phi(v) = i)$
(ii) Restrict the coloring ψ_{i-1} to $V(G_{i-1}) \setminus V(H_i)$
(iii) Extend the restriction to a coloring ψ_i of $G_i = G[V(G_{i-1}) \cup V(H_i)]$
- 6: **return** $\psi_{C^{2C}}$

D. Outline of Paper

In Section III, we prove Theorem I.2 and Theorem I.3, namely, by proving the correctness and runtime complexity of Algorithm 1. In Section IV, we prove Theorem II.4 and hence also its special case Theorem II.3. Finally in Section V, we prove Theorem II.7.

III. ALGORITHM ANALYSIS

In this section, we prove the correctness, termination and runtime complexity of Algorithm 1. First, we need the following lemma to lower bound the size of H in Step 2. Crucially, we need Theorem II.7 to do this.

Lemma III.1. Let Σ be a surface and let C be as in Theorem II.7 for Σ . Let G be a graph embeddable in Σ having girth at least $g \in \{3, 4, 5\}$ such that $v(G) \geq Cg(\Sigma)$. If S is the set of all vertices in G that are contained in an $(8 - g)$ -deletable C -pocket, then $|S| \geq \frac{v(G)}{2C}$.

Proof. Since $v(G) \geq Cg(\Sigma)$, we have by Theorem II.7 that there exists an $(8 - g)$ -deletable C -wallet \mathcal{H} of G . By definition then $|\{H \in \mathcal{H}\}| \geq \frac{1}{2C}v(G)$. Hence $\sum_{H \in \mathcal{H}} v(H) \geq \frac{1}{2C}v(G)$. Yet $\bigcup_{H \in \mathcal{H}} V(H) \subseteq S$. Since the subgraphs in \mathcal{H} are pairwise vertex-disjoint, it follows that $|S| \geq \frac{1}{2C}v(G)$. \square

Lemma III.1 implies an upper bound on G_0 in Step 2 as follows.

Corollary III.2. If G is as in Algorithm 1 and G_0 is defined as in Step 2, then $v(G_0) \leq (1 - \frac{1}{2C})v(G)$.

Proof. Since $V(H)$ is precisely the set of vertices contained in an $(8 - g)$ -deletable C -pocket, we find by Lemma III.1 that $v(H) \geq \frac{1}{2C}v(G)$. Hence $v(G_0) \leq v(G) - v(H) \leq (1 - \frac{1}{2C})v(G)$ as desired. \square

We are now ready to prove the termination and correctness of Algorithm 1.

Lemma III.3. *Algorithm 1 terminates and returns an L -coloring of G .*

Proof. We prove this by induction on $v(G)$. If $v(G) \leq Cg(\Sigma)$ as tested in Step 1, then the algorithm returns a coloring by exhaustive search. That search runs in constant time. Note that for planar graph an L -coloring always exists since L is a type 345 list-assignment (this follows from the results of Thomassen mentioned in the introduction; see Lemma V.1 for more details). As for graphs embedded in surfaces other than the plane, we assumed that G has an L -coloring and hence so does every subgraph of G .

So we may assume $v(G) \geq Cg(\Sigma)$. Step 2 is simply definitional for the purposes of terminations and correctness. By Corollary III.2, $v(G_0) \leq (1 - \frac{1}{2C})v(G) < v(G)$. Hence by induction, Step 3 returns an L -coloring ψ_0 of G_0 .

As for Step 4, we note that $\Delta(H) \leq C$ since every vertex in H has degree at most C in G . Since $C \geq 2$, it follows that $\Delta(H^{2C}) < C^{2C}$. Thus H^{2C} has a C^{2C} -coloring using the greedy bound of $\Delta(H^{2C}) + 1$.

As for Step 5, we now prove by induction that for each $i \in \{1, \dots, C^{2C}\}$ there exists an L -coloring ψ_i of G_i as in Step 5(iii) whose restriction to $V(G_{i-1}) \setminus V(H_i)$ agrees with ψ_{i-1} . Let $u \neq v \in V(H)$ such that $\phi(u) = \phi(v) = i$. Since $v(H_u), v(H_v) \leq C$ and the distance between u and v in G is at least $2C + 1$ as $\phi(u) = \phi(v)$, we have that H_u and H_v are non-touching. Thus the set of subgraphs $\mathcal{H} = \{H_v : \phi(v) = i\}$ are pairwise non-touching.

Now we define a list assignment L_i of H_i as follows. For each $v \in V(H_i)$, let $L_i(v) = L(v) \setminus \{\psi_{i-1}(u) : u \in N(v) \cap (V(G_{i-1}) \setminus V(H_i))\}$. Now $|L_i(v)| \geq |L(v)| - (d_G(v) - d_{H_i}(v)) = 8 - g - (d_G(v) - d_{H_i}(v))$. Let $H_v \in \mathcal{H}$. Since H_v is $(8 - g)$ -deletable in G , it follows by definition that there exists an L_i -coloring ψ_v of H_v . Now we define ψ_i as follows: let $\psi_i(w) = \psi_{i-1}(w)$ if $w \in V(G_{i-1}) \setminus V(H_i)$ and otherwise let $\psi_i(w) = \psi_v(w)$ where $w \in H_v$. By the definition of L_i and since the H_v are pairwise non-touching, it follows that ψ_i is an L -coloring of G_i as desired.

Hence in Step 6, Algorithm 1 returns $\psi_{C^{2C}}$ which is an L -coloring of $G_{C^{2C}} = G$ as desired. \square

Next we prove that Algorithm 1 runs in linear-time in the centralized setting as follows.

Lemma III.4. *Algorithm 1 runs in $O(v(G))$ -time.*

Proof. We proceed by induction on $n = V(G)$. We prove the runtime is an for some constant a as follows. Specifically $a = 2CB$ where B is as defined below.

We note that Step 1 runs in constant time. As for Step 2, we note that deciding if a given constant-size subgraph is r -

deletable can be done in constant time (since it is necessary to only check a finite number of list assignments). Thus for each vertex v , deciding if an H_v exists and finding one if it does can be done in constant time (since there are only a constant number of candidates as they have constant-size, have constant-degree in G , are connected and contain v). Hence Step 2 runs in $O(n)$ time (admittedly for some very large constant).

Finding G_0 for use in Step 3 can also be done in linear time (since planar graphs have at most $3n$ edges). Similarly for Step 4, finding H and H^{2C} can be done in linear time (since H has maximum degree at most C).

Step 4 then runs in $C^{2C}n$ time since $\Delta(H^{2C}) < C^{2C}$ and a $(\Delta(H) + 1)$ -coloring of H can be found in $(\Delta(H) + 1)n$ time by greedy coloring. As for Step 5, for each $i \in \{1, \dots, C^{2C}\}$, steps 5(ii) and 5(iii) can be done in linear time. Hence step 5 runs in linear-time.

Altogether each run of the algorithm (not counting recursion) is linear-time, specifically at most Bn for some universal constant B (depending only on C). By Corollary III.2, $v(G_0) \leq (1 - \frac{1}{2C})n$. Let $r = 1 - \frac{1}{2C}$. By induction, Step 3 takes $a|V(G_0)| = a(rn)$ time. Hence the total runtime of the algorithm is at most $Bn + arn = (B + ar)n$ which is at most an as desired if $B + ar \leq a$, that is, $a \geq \frac{B}{1-r}$. \square

Next we prove the runtime of Algorithm 1 in the deterministic distributed setting.

Lemma III.5. *Algorithm 1 runs in $O(\log v(G))$ -rounds in the deterministic distributed setting.*

Proof. We proceed by induction on $n = V(G)$.

Step 1 runs in $Cg(\Sigma)$ rounds (i.e. a constant number). Step 2 runs in C rounds (i.e. a constant number). For Step 4, we use the deterministic distributed algorithm of Fragniaud, Heinrich and Kosowski [11] to find a $(\Delta(H^{2C}) + 1)$ -coloring of H^{2C} in $\tilde{O}(\sqrt{\Delta}) + O(\log^* n)$ rounds. Since $\Delta(H) < C^{2C}$, this yields a C^{2C} -coloring ϕ of H^{2C} as desired. The number of rounds required for Step 3 is thus $\tilde{O}(C^C) + O(\log^* n) = O(\log^* n)$. As for Step 5, for each $i \in \{1, \dots, C^{2C}\}$, steps 5(i)-(iii) run in C rounds. Hence Step 5 runs in C^{2C+1} rounds (i.e. a constant number).

By Corollary III.2, $v(G_0) \leq (1 - \frac{1}{2C})n$. Thus the recursive step (Step 3) takes by induction $O(\log(v(G_0))) = O((1 - \frac{1}{2C})n)$ rounds. Since Step 4 is performed separately while the recursion runs, the total runtime is $O(\log^* n) + O(\log n) = O(\log n)$ rounds. \square

Finally, we have that Theorem I.2 and I.3 follow immediately from the above lemmas.

Proof of Theorem I.2. Follows from Lemmas III.3 and III.4. \square

Proof of Theorem I.3. Follows from Lemmas III.3 and III.5. \square

IV. SURFACE STRUCTURE THEOREM

Theorems II.3 and its generalization Theorem II.4 are mostly straightforward consequences of the following theorem of Lipton and Tarjan [14] about the separator hierarchy.

Lemma IV.1. *For every proper minor-closed family \mathcal{F} and every $\varepsilon > 0$, there exists a constant C such that the following holds: If G is a graph in \mathcal{F} , then there exists $X \subseteq V(G)$ with $|X| \leq \varepsilon v(G)$ such that every component of $G - X$ has size at most C .*

The proof of the result is essentially iteratively finding sublinear separators and taking X to be their union, wherein we terminate the procedure right before X becomes too large, whence the size of components in $G - X$ will be a small constant depending on ε .

We recall that for every proper minor-closed family \mathcal{F} , there exists $c_{\mathcal{F}}$ such that for every graph $G \in \mathcal{F}$, we have that $e(G) \leq c_{\mathcal{F}}v(G)$. That is, graphs in proper minor-closed families have a linear number of edges. Such a result was first shown by Mader [15], [16]. Using this, we obtain the following proposition.

Proposition IV.2. *For every proper minor-closed family \mathcal{F} and every $\varepsilon > 0$, there exists a constant C such that the following holds: If G is a graph in \mathcal{F} , then there exists $X \subseteq V(G)$ with $|X| \leq \varepsilon v(G)$ such that every $v \in V(G) \setminus X$ has degree at most C in G .*

Proof. Let $C = \frac{2c_{\mathcal{F}}}{\varepsilon}$. Let $X = \{v \in V(G) : d(v) > C\}$. Note that $2e(G) = \sum_{v \in V(G)} d(v) > |X|C$. Yet $e(G) \leq c_{\mathcal{F}}v(G)$ for some constant $c_{\mathcal{F}}$ as noted above. Hence $|X| \leq \frac{2e(G)}{C} \leq \frac{2}{C}(c_{\mathcal{F}}v(G)) = \varepsilon v(G)$ as desired. \square

Combining Lemma IV.1 with Proposition IV.2 yields the following corollary.

Corollary IV.3. *For every proper minor-closed family \mathcal{F} and every $\varepsilon > 0$, there exists a constant C such that the following holds: If G is a graph in \mathcal{F} , then there exists $X \subseteq V(G)$ with $|X| \leq \varepsilon v(G)$ such that every component of $G - X$ is a C -pocket of G .*

Proof. Let C_1 be as in Lemma IV.1 for \mathcal{F} and $\frac{\varepsilon}{2}$. By Theorem IV.1, there exists $X_1 \subseteq V(G)$ with $|X_1| \leq \frac{\varepsilon}{2}v(G)$ such that every component of $G - X_1$ has size at most C_1 .

Let C_2 be as in Proposition IV.2 for \mathcal{F} and $\frac{\varepsilon}{2}$. By Proposition IV.2, there exists $X_2 \subseteq V(G)$ with $|X_2| \leq \frac{\varepsilon}{2}v(G)$ such that every $v \in V(G) \setminus X_2$ has degree at most C_2 .

Let $C = \max\{C_1, C_2\}$. Let $X = X_1 \cup X_2$. Now $|X| \leq |X_1| + |X_2| \leq \varepsilon v(G)$. Moreover, every component of $G - X$ is a subgraph of a component of $G - X_1$ and hence has size at most $C_1 \leq C$. On the other hand, every vertex in $V(G) \setminus V(X)$ is a vertex in $V(G) \setminus V(X_2)$ and hence has degree at most $C_2 \leq C$ in G . Thus every component of $G - X$ is a C -pocket of G and hence X is as desired. \square

Next we prove a crucial lemma which shows how common purses are. To that end, recall the following definitions.

Definition IV.4. A vertex (resp. edge) amalgamation of two graphs G_1 and G_2 is a graph obtained by identifying a vertex (resp. edge) in G_1 with a vertex in G_2 .

Thus a vertex amalgamation is simply a 1-sum, while an edge amalgamation is a 2-sum where the edge is retained.

Definition IV.5. The (Euler) genus of graph G is the minimum k such that G embeds in a surface of Euler genus k .

We need the following result of Miller [17].

Theorem IV.6 (Miller [17]). *Euler genus is additive under vertex and edge amalgamations.*

For ease of reading in the proof, we adopt the following notation: if \mathcal{H} is a set of subgraphs of a graph G , we let $V(\mathcal{H}) = \bigcup_{H \in \mathcal{H}} V(H)$, let $v(\mathcal{H}) = |V(\mathcal{H})|$ and finally let $|\mathcal{H}|$ denote the number of subgraphs in \mathcal{H} .

Lemma IV.7. *If G is a connected graph embeddable in a surface Σ and $X \subseteq V(G)$, then there are most $12k|X| + 16kg(\Sigma)$ vertices in $G - X$ that are in components of $G - X$ that are neither purses of G nor k -deep in G .*

Proof. Let $\mathcal{H}_0 = \{H : H \text{ is a component of } G - X : H \text{ is neither a purse nor } k\text{-deep}\}$. Consider the following partition of \mathcal{H}_0 :

- $\mathcal{H}_1 = \{H \in \mathcal{H}_0 : \text{the coboundary of } H \text{ has size one}\}$,
- $\mathcal{H}_2 = \{H \in \mathcal{H}_0 : \text{the coboundary of } H \text{ has size two}\}$,
- $\mathcal{H}_3 = \{H \in \mathcal{H}_0 : \text{the coboundary of } H \text{ has size at least three}\}$.

We further partition \mathcal{H}_2 as follows:

- $\mathcal{H}_{2,1} = \{H \in \mathcal{H}_2 : \text{there does not exist } H' \neq H \in \mathcal{H}_2 \text{ with the same coboundary as } H\}$,
- $\mathcal{H}_{2,2} = \mathcal{H}_2 \setminus \mathcal{H}_{2,1}$.

Claim IV.8.

$$v(\mathcal{H}_1) + v(\mathcal{H}_{2,2}) \leq 4kg(\Sigma).$$

Proof. Let P be the set of unordered pairs $x \neq y \in V(G)$ such that there exists $H \in \mathcal{H}_{2,2}$ with coboundary $\{x, y\}$. For each $\{x, y\} \in P$, choose one $H_{xy} \in \mathcal{H}_2$ such that H_{xy} has coboundary $\{x, y\}$.

Let $G_1 = G[X \cup V(\mathcal{H}_1) \cup V(\mathcal{H}_{2,2})]$. Note that G_1 is a subgraph of G . Let G'_1 be obtained from G_1 by for each $\{x, y\} \in P$, deleting H_{xy} and adding the edge xy to G . Note that G'_1 is a minor of G_1 and hence is a minor of G . Thus G'_1 has Euler genus at most $g(\Sigma)$.

Let $H = G[X] + \bigcup_{\{x,y\} \in P} xy$. Now G'_1 is the vertex/edge amalgamations of H and $|\mathcal{H}_1| + |\mathcal{H}_{2,2}| - |P|$ non-planar graphs. Since G'_1 has Euler genus at most $g(\Sigma)$ and non-planar graphs have Euler genus at least one, it follows from Theorem IV.6 that

$$|\mathcal{H}_1| + |\mathcal{H}_{2,2}| - |P| \leq g(\Sigma).$$

Note that $2|P| \leq |\mathcal{H}_{2,2}|$ and hence $2|\mathcal{H}_1| + |\mathcal{H}_{2,2}| \leq 2g(\Sigma)$.

Let $H_1 \in \mathcal{H}_1$. Note that H_1 has a coboundary of size one. Yet H_1 is not k -deep. Hence $|H_1| \leq k$. Thus $v(\mathcal{H}_1) \leq k|\mathcal{H}_1|$.

Similarly, we find that $v(\mathcal{H}_{2,2}) \leq 2k|\mathcal{H}_{2,2}|$ and the claim follows. \square

Claim IV.9.

$$v(\mathcal{H}_{2,1}) \leq 6k(|X| + g(\Sigma)).$$

Proof. Let P' be the set of unordered pairs $x \neq y \in V(G)$ such that there exists $H \in \mathcal{H}_{2,1}$ with coboundary $\{x, y\}$. Let $G_2 = X + \bigcup_{\{x,y\} \in P'} xy$. Note that G_2 is a minor of G and hence has Euler genus at most $g(\Sigma)$. By Euler's formula, $e(G_2) \leq 3(|X| + g(\Sigma))$. Hence,

$$|\mathcal{H}_{2,1}| = e(G_2) \leq 3(|X| + g(\Sigma)).$$

Since every element of $\mathcal{H}_{2,1}$ has a coboundary of size two and is not k -deep, we have that $v(\mathcal{H}_{2,1}) \leq 2k|\mathcal{H}_{2,1}|$ and the claim follows. \square

Claim IV.10.

$$v(\mathcal{H}_3) \leq 6k(|X| + g(\Sigma)).$$

Proof. Let G_3 be obtained from $G[X \cup V(\mathcal{H}_3)]$ by contracting each $H \in \mathcal{H}_3$ to a new vertex v_H and deleting loops and parallel edges. Now G_3 is a minor of G and hence has Euler genus at most $g(\Sigma)$. Note that G_3 is bipartite. By Euler's formula, $e(G_3) \leq 2(v(G_3) + g(\Sigma))$. Yet every vertex in $V(G_3) \setminus X$ has degree at least three. Hence $e(G_3) \geq 3|\mathcal{H}_3|$. Since $v(G_3) = |X| + |\mathcal{H}_3|$, we find that

$$|\mathcal{H}_3| \leq 2(|X| + g(\Sigma)),$$

and hence

$$e(G_3) \leq 6(|X| + g(\Sigma)).$$

Let $H \in \mathcal{H}_3$. Note that v_H has degree in G_3 equal to the size of the coboundary of H in G . Moreover, since H is not k -deep, it follows that $v(H) \leq k \cdot \deg_{G_3}(v_H)$. Since G_3 is bipartite, it follows that $e(G_3) = \sum_{H \in \mathcal{H}_3} \deg_{G_3}(v_H)$ and hence

$$v(\mathcal{H}_3) \leq k \cdot e(G_3) \leq 6k(|X| + g(\Sigma)),$$

as claimed. \square

Combining Claims IV.8, IV.9 and IV.10, we find that

$$v(\mathcal{H}_0) \leq 12k|X| + 16kg(\Sigma),$$

as desired. \square

We are now ready to prove Theorem II.4 of which Theorem II.3 is a special case.

Proof of Theorem II.4. Let \mathcal{F} be the family of graphs embeddable in Σ . Recall that \mathcal{F} is a minor-closed family. Let C_0 be as C in Corollary IV.3 for \mathcal{F} with $\varepsilon = \frac{1}{58k}$. Since $k \geq 1$, it follows that $\varepsilon < 1$ and hence $C_0 \geq 1$. Let $C = 58kC_0$. Since $v(G) \geq Cg(\Sigma)$ by assumption, we have that $g(\Sigma) \leq \frac{1}{58kC_0}v(G)$ which is at most $\frac{1}{58k}v(G)$ since $C_0 \geq 1$.

By Corollary IV.3, there exists $X \subseteq V(G)$ with $|X| \leq \varepsilon v(G)$ such that every component of $G - X$ is a C_0 -pocket. Let $\mathcal{H}_1 = \{H : H \text{ is a component of } G - X, H \text{ is neither a purse nor } k\text{-deep}\}$. By Lemma IV.7, we have that

$$v(\mathcal{H}_1) \leq 12k|X| + 16kg(\Sigma).$$

Let $\mathcal{H} = \{H \text{ is a component of } G - X : H \notin \mathcal{H}_1\}$. Now

$$v(\mathcal{H}) \geq v(G) - |X| - v(\mathcal{H}_1) \geq v(G) - (12k+1)|X| - 16kg(\Sigma).$$

Since $k \geq 1$, we have that $12k + 1 \leq 13k$. Since $|X| \leq \varepsilon v(G) = \frac{1}{58k}v(G)$ and $g(\Sigma) \leq \frac{1}{58k}v(G)$, we find that

$$v(\mathcal{H}) \geq v(G) - (13k + 16k) \cdot \frac{v(G)}{58k} \geq \frac{1}{2}v(G).$$

Since every $H \in \mathcal{H}_2$ has at most C vertices, it follows that

$$|\mathcal{H}| \geq \frac{1}{2C}v(G),$$

and hence \mathcal{H} is a k -deep C -wallet as desired. \square

V. LINEAR SIZE FOR NO DELETABLE SUBGRAPHS

We prove Theorem II.7 via Theorem II.4. Namely Theorem II.4 guarantees us a k -deep C -wallet, that is a set of linearly many pairwise non-touching subgraphs, each of which is either a C -purse or a k -deep C -pocket. Thus to prove Theorem II.7, it suffices to prove the following two lemmas.

Lemma V.1. *Let $g \in \{3, 4, 5\}$. If D is a purse of a graph G having girth at least g , then D is $(8 - g)$ -deletable in G .*

Lemma V.2. *For each surface Σ , there exists $k > 0$ such that the following holds: Let $g \in \{3, 4, 5\}$ and let G be a graph embeddable in Σ having girth at least g . If D is an induced k -deep subgraph of G , then there exists $X \subseteq V(D)$ such that $G[X]$ is $(8 - g)$ -deletable in G .*

In Section 5.1, we establish Lemma V.1 using the standard results of Thomassen on list-coloring planar graphs.

On the other hand, the proof of Lemma V.2 will require the use of the theory of hyperbolic families as well as the deep results of the author and Thomas [23] on the strong hyperbolicity of graphs that are critical for 5-list-coloring, and of the author [19] on the strong hyperbolicity of graphs of girth five that critical for 3-list-coloring. We use these and other results to show that embedded graphs having girth at least g without $(8 - g)$ -deletable subgraphs form a *strongly hyperbolic family* (see Definition V.7 below). The Strongly Hyperbolic Structure Theorem (Theorem 7.2 in [23]) will then imply that for every surface there exists a k such that every k -deep subgraph of a graph embedded in that surface contains an $(8 - g)$ -deletable subgraph, thereby completing the proof of Theorem II.7.

We are now ready to prove Theorem II.7 assuming Lemmas V.1 and V.2.

Proof of Theorem II.7. Fix a surface Σ . Let k be as in Lemma V.2 for Σ . Let C be as in Theorem II.4 for Σ and k . Thus by Theorem II.4, there exists a k -deep C -wallet \mathcal{H} of G . Recall that a C -wallet is a set of at least $\frac{1}{2C}v(G)$ pairwise non-touching C -pockets. Recall that a C -wallet is k -deep if for every $D \in \mathcal{H}$, either D is a purse or k -deep.

For each $D \in \mathcal{H}$ such that D is a purse, we have by Lemma V.1 that D is $(8-g)$ -deletable in G . For each $D \in \mathcal{H}$ such that D is k -deep, we have by Lemma V.2 that there exists $X_D \subseteq V(D)$ such that $G[X_D]$ is $(8-g)$ -deletable.

Let $\mathcal{H}' = \{D \in \mathcal{H} : D \text{ is a purse}\} \cup \{X_D : D \in \mathcal{H} \text{ is } k\text{-deep}\}$. Thus \mathcal{H}' is a set of at least $\frac{1}{2C}v(G)$ pairwise non-touching subgraphs of G , each of which is $(8-g)$ -deletable, as desired. \square

Outline of Section. In Section V-A, we prove Lemma V.1. In Section V-B, we recall the many definitions and theorems of hyperbolic families that we need for the proof of Lemma V.2. In Section V-C, we define the crucial object of study: a certain family of embedded graphs with no deletable subgraphs and then provide a proof of Lemma V.2 assuming that family is strongly hyperbolic. In Section V-D, we collect the other previous results about hyperbolic families that we need for the proof. In Section V-E, we prove that our crucial family is hyperbolic (see Definition V.6; this is a weaker notion than strongly hyperbolic). Finally in Section V-F, we then use the result of Section V-E to prove that the crucial family is in fact strongly hyperbolic.

A. Deletability for Purses

In order to prove Theorem II.7 via Theorem II.4, we will need to prove that every C -purse contains an $(8-g)$ -deletable subgraph. In fact, we will prove the stronger theorem that every C -purse is itself $(8-g)$ -deletable. This mostly follows from theorems of Thomassen as we will show. First we prove for planar graphs that every L -coloring of a short precolored path extends to an L -coloring of the whole graph using results of Thomassen as follows.

Theorem V.3. *Let G be a planar graph having girth at least $g \in \{3, 4, 5\}$ and let L be a type 345 list-assignment for G . If P is a path in G such that $v(P) \leq g-1$, then every L -coloring of P extends to an L -coloring of G .*

Proof. For $g = 3$, this follows from Thomassen's stronger inductive theorem about 5-list-coloring planar graphs (specifically the main Theorem in [25]). For $g = 5$, this follows from Thomassen's stronger inductive theorem about 3-list-coloring planar graphs of girth at least five (specifically Theorem 2.1 in [29], the main result of that paper).

For $g = 4$, this is rather straightforward as follows. We proceed by induction on $v(G)$. We may assume $v(G) \geq 3$ and that $V(P) \neq V(G)$ as otherwise the result is trivial. We may assume by induction that G is connected.

We claim there exists a vertex $v \in V(G) \setminus V(P)$ of degree at most 3. Suppose not. By Euler's formula since $v(G) \geq 3$ and G is triangle-free, we have that $e(G) \leq 2v(G) - 4$.

Thus, $\sum_{v \in V(G)} d(v) \leq 4v(G) - 8$. Yet $\sum_{v \in V(P)} d(v) \geq 4(v(G) - v(P))$. Combining the previous inequalities, we find that $\sum_{v \in V(P)} d_G(v) \leq 4v(P) - 8$. But $\sum_{v \in V(P)} d_G(v) \geq \sum_{v \in V(P)} d_P(v) = 2v(P) - 2$ since P is a path. It follows that $v(P) = 3$ and P is a component of G , contradicting that G is connected. This proves the claim.

By induction, ϕ extends to an L -coloring of $G-v$ and hence to an L -coloring of G as desired. \square

We now use Theorem V.3 to prove Lemma V.1 as follows.

Proof of Lemma V.1. Suppose not. That is, there exists a list-assignment L_0 such that $|L_0(v)| \geq 8-g - (d_G(v) - d_D(v))$ for each $v \in V(D)$ and D is not L_0 -colorable.

Let S be the coboundary of D in G . First suppose $S = \emptyset$. Then L_0 is an $(8-g)$ -list-assignment of D . By Theorem V.3 with $P = \emptyset$, there exists an L_0 -coloring of D , a contradiction.

Next suppose $|S| = 1$. Let $S = \{x\}$. Let $G' = G[V(D) \cup \{x\}]$. Let c be a new color (that is $c \notin \bigcup_{v \in V(D)} L_0(v)$). Define a new list assignment L of G' as follows. Let $L(x) = \{c\} \cup R$ where R is a set of $7-g$ arbitrary colors. For each $v \in V(D)$, let $L(v) = L_0(v) \cup \{c\}$.

Now L is an $(8-g)$ -list-assignment of G' . It follows from Theorem V.3 with $P = x$ that there exists an L -coloring ϕ of G' such that $\phi(x) = c$. But then ϕ induces an L_0 -coloring of D , a contradiction.

Finally suppose $|S| = 2$. Let $S = \{x, y\}$. Let $G' = G[V(D) \cup \{x, y\}]$. Let G'' be obtained from G' as follows: delete the edge xy if it exists and add a path $P = xv_1 \dots v_{g-3}y$ with new vertices v_1, \dots, v_{g-3} . Now G'' has girth at least g . Since D is a purse, it follows that G'' is planar.

Let ϕ be a coloring of P with entirely new colors (i.e. $\phi(u) \notin \bigcup_{v \in V(H)} L_0(v)$ for every $u \in V(P)$) such that $\phi(x) \neq \phi(y)$. Define a new list assignment L of G' as follows. For each $u \in V(P)$, let $L(u) = \{\phi(u)\} \cup R$ where R is a set of $7-g$ arbitrary colors. For each $v \in V(D)$, let $L(v) = L_0(v) \cup \{\phi(u) : u \in N(v) \cap V(P)\}$. Now L is an $(8-g)$ -list-assignment of G'' . It follows from Theorem V.3 that ϕ extends to an L -coloring of G'' . But then this induces an L_0 -coloring of D , a contradiction. \square

B. Theory of Hyperbolic Families

We will need to recall a number of definitions from [23]. First we recall the definition of a graph with rings and how they embed in a surface.

Definition V.4 (Graph with Rings - Definition 3.1 in [23]). A *ring* is a cycle or a complete graph on one or two vertices. A *graph with rings* is a pair (G, \mathcal{R}) , where G is a graph and \mathcal{R} is a set of vertex-disjoint rings in G .

Definition V.5 (Embedding Graphs with Rings - Definition 3.2 in [23]). We say that a graph G with rings \mathcal{R} is *embedded in a surface* Σ if the underlying graph G is embedded in Σ in such a way that for every ring $R \in \mathcal{R}$ there exists a component Γ of the boundary of Σ such that R is embedded in Γ , no other

vertex or edge of G is embedded in Γ , and every component of the boundary of Σ includes some ring of G .

Now let us state the formal definition of a hyperbolic family. Informally the definition says that for every graph in the family the number of vertices inside a disk is at most linear in the number of vertices on the boundary of that disk.

Definition V.6 (Hyperbolic Family - Definition 5.1 in [23]). Let \mathcal{F} be a family of non-null embedded graphs with rings. We say that \mathcal{F} is *hyperbolic* if there exists a constant $c > 0$ such that if $G \in \mathcal{F}$ is a graph with rings that is embedded in a surface Σ , then for every closed curve $\gamma : S^1 \rightarrow \Sigma$ that bounds an open disk Δ and intersects G only in vertices, if Δ includes a vertex of G , then the number of vertices of G in Δ is at most $c(|\{x \in S^1 : \gamma(x) \in V(G)\}| - 1)$. We say that c is a *Cheeger constant* of \mathcal{F} .

Finally let us also state the definition of a strongly hyperbolic family. Informally the definition extends the linearity property of hyperbolic families from disks to also include annuli (a.k.a. cylinders).

Definition V.7 (Strongly Hyperbolic Family - Definition 7.1 in [23]). Let \mathcal{F} be a hyperbolic family of embedded graphs with rings, let c be a Cheeger constant for \mathcal{F} , and let $d := \lceil 3(2c+1) \log_2(8c+4) \rceil$. We say that \mathcal{F} is *strongly hyperbolic* if there exists a constant c_2 such that for every $G \in \mathcal{F}$ embedded in a surface Σ with rings and for every two disjoint cycles C_1, C_2 of length at most $2d$ in G , if there exists a cylinder $\Lambda \subseteq \Sigma$ with boundary components C_1 and C_2 , then Λ includes at most c_2 vertices of G . We say that c_2 is a *strong hyperbolic constant* for \mathcal{F} .

A key result we need from [23] is that the number of vertices in a graph in a strongly hyperbolic family is linear in the sum of its genus and ring vertices as follows. Before stating that result, we need one more technical but rather innocuous definition.

Definition V.8 (Closed under curve cutting - appears after Theorem 1.1 in [23]). A family \mathcal{F} of embedded graphs is *closed under curve cutting* if for every embedded graph $G \in \mathcal{F}$ embedded in a surface Σ and every simple closed curve ξ in Σ whose image is disjoint from G , if Σ' denotes the surface obtained from Σ by cutting open along ξ and attaching disk(s) to the resulting curve(s), then the embedded graph G in Σ' belongs to \mathcal{F} .

Theorem V.9 (A Simplified Form of Theorem 7.2 in [23]). *Let \mathcal{F} be a strongly hyperbolic family of embedded graphs with rings such that \mathcal{F} is closed under curve-cutting. Then there exists a constant $c_{\mathcal{F}}$ such the the following holds: if $G \in \mathcal{F}$ is embedded in a surface Σ of Euler genus g with a total of R ring vertices, then $v(G) \leq c_{\mathcal{F}}(g + R)$.*

C. A Family with No Deletable Subgraphs

Next we define our crucial family of embedded graphs with rings that we need to prove Lemma V.2.

Definition V.10. For each $g \in \{3, 4, 5\}$, let \mathcal{F}_g be the family of embedded graphs with rings (G, \mathcal{R}) such that $G - \bigcup_{R \in \mathcal{R}} E(R)$ has girth at least g and there does not exist an $(8-g)$ -deletable subgraph of G whose vertices all lie in $G - \bigcup_{R \in \mathcal{R}} V(R)$. Let $\mathcal{F}_{345} = \mathcal{F}_3 \cup \mathcal{F}_4 \cup \mathcal{F}_5$.

We note that \mathcal{F}_{345} is clearly closed under curve cutting. Using results by the author and Thomas, we will prove in Section V-F the following theorem.

Theorem V.11. *\mathcal{F}_{345} is strongly hyperbolic.*

The following is an immediate corollary of Theorem V.11 and Theorem V.9.

Corollary V.12. *There exists $k > 0$ such that the following holds: Let G be embedded in a surface Σ of genus $g(\Sigma)$ having girth at least $g \in \{3, 4, 5\}$. If H is a subgraph of G such that there does not exist an induced subgraph of $G - V(H)$ that is $(8-g)$ -deletable in G , then $v(G) \leq k(v(H) + g(\Sigma))$.*

Proof. By Theorem V.11, we have that \mathcal{F}_{345} is strongly hyperbolic. We let $k = c_{\mathcal{F}_{345}}$ be as in Theorem V.9.

Now let $G' = G - E(H)$. We view G' as a graph with rings where each vertex of H is a ring. Now G' also has girth at least g and is embedded in Σ . Since there does not exist a subgraph of $G - V(H)$ that is $(8-g)$ -deletable in G , it follows that $G' \in \mathcal{F}_{345}$. Hence by Theorem V.9, we have that $v(G) \leq k(v(H) + g(\Sigma))$ as desired. \square

We now prove Lemma V.2 assuming Theorem V.11.

Proof of Lemma V.2. Let k_0 be as in Corollary V.12. We let $k = k_0(g(\Sigma) + 1) + 1$. Suppose for a contradiction that there does not exist $X \subseteq V(D)$ such that $G[X]$ is $(8-g)$ -deletable in G .

Let H be the coboundary of D in G . Let $G' = G[V(H) \cup V(D)]$. Note then that there does not exist an induced subgraph of $G' - V(H)$ that is deletable in G' (as otherwise it would be $(8-g)$ -deletable in G). Now by Corollary V.12, we have that $v(G') \leq k_0(v(H) + g(\Sigma))$.

Since D is k -deep, $v(H)$ is non-empty by definition and hence $v(H) \geq 1$. Thus $k_0(v(H) + g(\Sigma)) < kv(H)$ and so $v(G') < kv(H)$. Rewriting, we have that $v(H) > \frac{1}{k}v(G') > \frac{1}{k}v(D)$. Yet since D is k -deep, we have by definition that $v(H) \leq \frac{1}{k}v(D)$, a contradiction. \square

D. Previous Results on Hyperbolic Families

In order to prove Theorem V.11, we will first prove that \mathcal{F}_{345} is hyperbolic and then use that in the proof that \mathcal{F}_{345} is strongly hyperbolic. However, we will need a number of previous results for both those proofs which we now collect.

First we recall the definition for being critical with respect to a subgraph for list-coloring.

Definition V.13. Let G be a graph and let H be a subgraph of G . If L is a list assignment of G , then we say G is *H -critical with respect to L* if the following holds: for every proper subgraph G' of G containing H , there exists an L -coloring of H that extends to G' but not to G . We say G is

H -critical for k -list-coloring if there exists a k -list-assignment such that G is H -critical with respect to L .

Definition V.14. For each $g \in \{3, 4, 5\}$, let \mathcal{G}_g be the family of embedded graphs with rings (G, \mathcal{R}) such that every cycle in G of length at most $(g - 1)$ is not null-homotopic and G is $\bigcup_{R \in \mathcal{R}} R$ -critical for $(8 - g)$ -list-coloring. Let $\mathcal{G}_{345} = \mathcal{G}_3 \cup \mathcal{G}_4 \cup \mathcal{G}_5$.

It was proved by the author and Thomas in [21] that \mathcal{G}_3 is hyperbolic and by Dvořák and Kawarabayashi in [8] that \mathcal{G}_5 is hyperbolic (the proof that \mathcal{G}_4 is hyperbolic is easy by comparison and can be found in [23]). The author and Thomas proved that \mathcal{G}_{345} is strongly hyperbolic as follows.

Theorem V.15. \mathcal{G}_{345} is strongly hyperbolic.

Proof. \mathcal{G}_3 is strongly hyperbolic by Lemma 7.5 in [23]. \mathcal{G}_4 is strongly hyperbolic by Lemma 7.6 in [23]. \mathcal{G}_5 is strongly hyperbolic by Lemma 7.8 in [23]. Hence \mathcal{G}_{345} is also strongly hyperbolic. \square

We remark that the proof of Lemmas 7.8 in [23] relies on the work of the author in [19] for 3-list-coloring girth five graphs. Similarly the proof of Lemma 7.5 relies on the work by the author and Thomas in the series of papers [20], [21], [22] (see [18] for a full proof). On the other hand, the proof of Lemma 7.6 for 4-list-coloring triangle-free graphs is once again quite easy by comparison (and can be found in Lemmas 5.10 and 7.6 in [23]).

Combining Theorem V.15 with Theorem V.9, we obtain the following corollary.

Corollary V.16. *There exists $k' > 0$ such that the following holds: Let G be embedded in a surface Σ of genus $g(\Sigma)$ having girth at least $g \in \{3, 4, 5\}$. If H is a subgraph of G such that G is H -critical for $(8 - g)$ -list-coloring, then $v(G) \leq k'(v(H) + g(\Sigma))$.*

Proof. By Theorem V.15, we have that \mathcal{G}_{345} is strongly hyperbolic. We let $k' = c_{\mathcal{G}_{345}}$ be as in Theorem V.9.

Now let $G' = G - E(H)$. We view G' as a graph with rings where each vertex of H is a ring. Now G' also has girth at least g and is embedded in Σ . Since G is H -critical for $(8 - g)$ -list-coloring, it follows that G' is also H -critical for $(8 - g)$ -list-coloring. Thus $G' \in \mathcal{G}_{345}$. Hence by Theorem V.9, we have that $v(G) \leq k'(v(H) + g(\Sigma))$ as desired. \square

To prove Theorem V.11, we will actually need a stronger result, namely instead of the hyperbolicity of \mathcal{G}_{345} , we will need a stronger density result (Theorem V.20 below). Theorem V.20 will be the key to proving that \mathcal{F}_{345} is hyperbolic. Finally to prove that \mathcal{F}_{345} is strongly hyperbolic, we will also use Theorem V.20 along with the strong hyperbolicity of \mathcal{G}_{345} .

To state the density result, we first define a crucial notion of the density of a graph over a subgraph.

Definition V.17. Let $g \in \{3, 4, 5\}$ and $\varepsilon > 0$. If G is a graph and H is a subgraph of G , define

$$d_{g,\varepsilon}(G|H) = (g - 2)(e(G) - e(H)) - (g + \varepsilon)(v(G) - v(H)).$$

We will also need the following easy proposition about $d_{g,\varepsilon}$.

Proposition V.18. *Let $g \in \{3, 4, 5\}$ and $\varepsilon > 0$. If $H \subseteq G' \subseteq G$, then*

$$d_{g,\varepsilon}(G|H) = d_{g,\varepsilon}(G|G') + d_{g,\varepsilon}(G'|H).$$

Proof. This follows from the definition of $d_{g,\varepsilon}$ since $e(G) - e(H) = (e(G) - e(G')) + (e(G') - e(H))$ and $v(G) - v(H) = (v(G) - v(G')) + (v(G') - v(H))$. \square

For a planar graph G of girth at least g , satisfying $d_{g,\varepsilon}(G|H) \geq 0$ implies an upper bound on $v(G)$ that is linear in $v(H)$ as our next lemma shows (a useful fact for proving the linear bounds required for hyperbolicity and strong hyperbolicity and hence the reason for the introduction of $d_{g,\varepsilon}$).

Lemma V.19. *Let $g \in \{3, 4, 5\}$ and $\varepsilon > 0$. If G is a planar graph having girth at least g and H is a subgraph of G such that $d_{g,\varepsilon}(G|H) \geq 0$, then*

$$v(G) \leq \frac{g + \varepsilon}{\varepsilon} v(H).$$

Proof. We have that $d_{g,\varepsilon}(G|H) = (g - 2)(e(G) - e(H)) - (g + \varepsilon)(v(G) - v(H)) \geq 0$. Rearranging, we have that

$$\varepsilon v(G) \leq (g + \varepsilon)v(H) - (g - 2)e(H) + ((g - 2)e(G) - gv(G)).$$

Since G is a planar graph having girth at least g , we have by Euler's formula that $e(H) \leq \frac{g}{g-2}v(G)$. Combining this observation with the fact that $e(H) \geq 0$, we find that

$$v(G) \leq \frac{(g + \varepsilon)}{\varepsilon} v(H),$$

as desired. \square

We may now state the density result which combines earlier theorems of the author [19] and the author and Thomas [23].

Theorem V.20. *There exists $\varepsilon > 0$ such that following holds: If G is a plane graph having girth at least $g \in \{3, 4, 5\}$ and H is a connected subgraph of G such that G is H -critical for $(8 - g)$ -list-coloring, then*

$$d_{g,\varepsilon}(G|H) \geq 0.$$

Proof. For $g = 3$, this is equivalent to Theorem 4.6 of the author and Thomas in [21], which constitutes the main result of that paper. For $g = 4$, the proof is straightforward and is implicit in Lemma 5.10 of the author and Thomas in [23]. For $g = 5$, this is a special case of Theorem 3.9 of the author in [19], which again constitutes the main result of that paper. \square

We note that combined with Lemma V.19, Theorem V.20 implies that the family \mathcal{G}_{345} is hyperbolic. This is because essentially hyperbolicity would be equivalent to showing that

$v(G)$ is linear in $v(H)$ for connected subgraphs H , while instead Theorem V.20 proves a stronger relation between the edges and vertices of G and H . Indeed these stronger statements were necessary for the inductive proofs of the hyperbolicity of \mathcal{G}_3 and \mathcal{G}_5 .

We note that the assumption that H is connected is related to hypercolity (being inside a disc) while the case of H having two connected components is related to strong hyperbolicity (being inside an annulus).

E. Hyperbolicity for No Deletable Subgraphs

We will use Theorem V.20, a statement stronger than the hyperbolicity of \mathcal{G}_{345} , to prove that the family \mathcal{F}_{345} is hyperbolic. First, we need the following easy proposition.

Proposition V.21. *Let G be a graph and H a proper subgraph of G such that $V(H) \neq V(G)$. If $G - V(H)$ is not r -deletable in G , then there exists a subgraph G' of G containing H such that G' is H -critical for r -list-coloring.*

Proof. Since $G - V(H)$ is not r -deletable in G , there exists a list-assignment L_0 such that $|L_0(v)| \geq r - d_H(v)$ for each $v \in V(G) \setminus V(H)$ and $G - V(H)$ is not L_0 -colorable.

Let $S = \bigcup_{u \in V(G) \setminus V(H)} L_0(u)$. Let $S' = \{c_v : v \in V(H)\}$ be a set of pairwise distinct new colors (that is $c_v \notin S$ for each $v \in V(H)$). Let R be an arbitrary set of $r - 1$ colors disjoint from S' .

Define a new list assignment L of G as follows. For each $v \in V(H)$, let $L(v) = \{c_v\} \cup R$. For each $u \in V(G) \setminus V(H)$, let $L(u) = L_0(u) \cup \{c_v : v \in N(u) \cap V(H)\}$. Now L is an r -list-assignment of G . Let ϕ be the coloring of H given by $\phi(v) = c_v$ for every $v \in V(H)$.

Since $G - V(H)$ is not L_0 -colorable, it follows that ϕ does not extend to an L -coloring of G . Let G' be an inclusion-wise minimal subgraph of G containing H such that ϕ does not extend to an L -coloring of G' . By the minimality of G' , ϕ extends to an L -coloring of every proper subgraph of G' containing H . Thus G' is H -critical with respect to L . Hence by definition, G' is H -critical for r -list-coloring, as desired. \square

Using Theorem V.20, we are now ready to prove the hyperbolicity of \mathcal{F}_{345} as follows. First we prove a density result as follows.

Lemma V.22. *Let ε be as in Theorem V.20. If G is a plane graph having girth at least $g \in \{3, 4, 5\}$ and H is a connected subgraph of G such that there does not exist $X \subseteq V(G) \setminus V(H)$ such that $G[X]$ is $(8 - g)$ -deletable in G , then*

$$d_{g,\varepsilon}(G|H) \geq 0.$$

Proof. We proceed by induction on $v(G) - v(H) + e(G) - e(H)$.

If $V(H) = V(G)$, then $d_{g,\varepsilon}(G|H) \geq 0$ as desired. So we may assume that $V(H) \neq V(G)$. By assumption, $G - V(H)$ is not $(8 - g)$ -deletable in G . By Proposition V.21, it follows that there exists a subgraph G' of G containing H such that G' is H -critical for $(8 - g)$ -list-coloring. Note that H is a proper

subgraph of G' by definition of H -critical. By Theorem V.20, we have that $d_{g,\varepsilon}(G'|H) \geq 0$. By Lemma V.1, it follows that G' is connected. Note that $v(G) - v(G') + e(G) - e(G') < v(G) - v(H) + e(G) - e(H)$. Hence by induction, $d_{g,\varepsilon}(G|G') \geq 0$. By Proposition V.18, we have that

$$d_{g,\varepsilon}(G|H) = d_{g,\varepsilon}(G|G') + d_{g,\varepsilon}(G'|H) \geq 0 + 0 = 0,$$

as desired. \square

Theorem V.23. *The family \mathcal{F}_{345} is hyperbolic.*

Proof. Let $c = \frac{6(5+\varepsilon)}{\varepsilon}$ where ε is as in Theorem V.20. We prove that c is a Cheeger constant for \mathcal{F}_{345} as follows.

Let G be a graph with rings \mathcal{R} embedded in a surface Σ of Euler genus $g(\Sigma)$ such that $G \in \mathcal{F}_{345}$, let R be the total number of ring vertices, and let $\xi : S_1 \rightarrow \Sigma$ be a closed curve that bounds an open disk Δ and intersects G only in vertices. To avoid notational complications we will assume that ξ is a simple curve; otherwise we split vertices that ξ visits more than once to reduce to this case. We may assume that Δ includes at least one vertex of G , for otherwise there is nothing to show. Let X be the set of vertices of G intersected by ξ .

Let G_0 be the subgraph of G consisting of all vertices and edges drawn in the closure of Δ . We now regard G_0 as a graph embedded in the closure of Δ inside the plane. We define a planar graph G_1 obtained from G_0 as follows: we add a vertex v in the complement of the closure of Δ ; for each vertex $x \in X$, we add the edge vx ; we then subdivide each edge incident with v exactly once.

Note that G_1 is planar and has girth at least g . Let $H = G_1[\{x\} \cup N_{G_1}(x) \cup X]$. Note that $G_0 - X$ is not $(8 - g)$ -deletable and hence $G_1 - V(H)$ is not $(8 - g)$ -deletable. It follows from Lemma V.1 that $G_1 - V(H)$ is not a purse. Since G_1 is planar, it thus follows that $v(H) \geq 2$.

Since $v(H) \geq 2$, we have $v(H) \leq 3|X| \leq 6(|X| - 1)$. Note that H is connected. Hence by Lemma V.22, $d_{g,\varepsilon}(G_1|H) \geq 0$. Thus by Lemma V.19, $v(G_1) \leq \frac{g+\varepsilon}{\varepsilon}v(H)$. But then

$$\begin{aligned} v(G_0) &\leq v(G_1) \leq \frac{g+\varepsilon}{\varepsilon}v(H) \\ &\leq \frac{6(5+\varepsilon)}{\varepsilon}(v(H) - 1) = c(v(H) - 1), \end{aligned}$$

as desired. \square

F. Strong Hyperbolicity for No Deletable Subgraphs

We now prepare to prove the strong hyperbolicity of \mathcal{F}_{345} . First we prove the following lemma.

Lemma V.24. *There exists $k > 0$ such that the following holds: If G is a connected planar graph having girth at least $g \in \{3, 4, 5\}$ and H_1, H_2 are non-empty vertex-disjoint connected subgraphs of G such that there does not exist $X \subseteq V(G) \setminus (V(H_1) \cup V(H_2))$ such that $G[X]$ is $(8 - g)$ -deletable in G , then the distance between H_1 and H_2 in G is at most $k(v(H_1) + v(H_2))$.*

Proof. We prove that $k = k' \frac{5+\varepsilon}{\varepsilon}$ suffices where k' is as in Corollary V.16.

Suppose not. Let $H = H_1 \cup H_2$. Let G' be a subgraph of G containing H such that every component of G' contains a vertex of H and $d_{g,\varepsilon}(G'|H) \geq 0$ and subject to that, $e(G') + v(G')$ is maximum.

Since G is planar, we have that G' is planar. Since $d_{g,\varepsilon}(G'|H) \geq 0$ and G' is planar, we have by Lemma V.19 that $v(G') \leq \frac{g+\varepsilon}{\varepsilon}v(H)$. Note that it also follows from Lemma V.19 that every component of G' contains a vertex in H and hence G' has at most two components.

First suppose that $V(G') = V(G)$. Then $v(G) = v(G') \leq \frac{g+\varepsilon}{\varepsilon}v(H)$. Since G is connected, it follows that the distance between H_1 and H_2 in G is at most $v(G) \leq \frac{5+\varepsilon}{\varepsilon}v(H) \leq k'(v(H_1) + v(H_2))$, a contradiction.

So we may assume that $V(G') \neq V(G)$. Since $G - V(G')$ is not $(8-g)$ -deletable in G' , we have by Proposition V.21 that there exists a subgraph G'' of G containing G' such that G'' is G' -critical for $(8-g)$ -list-coloring.

Now suppose that G' is connected. By Lemma V.20, $d_{g,\varepsilon}(G''|G') \geq 0$. By Proposition V.18,

$$d_{g,\varepsilon}(G''|H) = d_{g,\varepsilon}(G''|G') + d_{g,\varepsilon}(G'|H) \geq 0 + 0 \geq 0,$$

contradicting the maximality of G' .

So we may assume that G' is not connected. As noted above, we have that G' has exactly two components G_1 and G_2 . We may assume without loss of generality that $H_1 \subseteq G_1$ and $H_2 \subseteq G_2$.

Next suppose that G'' is not connected. Since G'' is planar and G'' is G' -critical for $(8-g)$ -list-coloring, it follows from Lemma V.1 that every component of G'' contains a vertex of G' . Hence G'' has two components G'_1 and G'_2 . We may assume without loss of generality that $G_1 \subseteq G'_1$ and $G_2 \subseteq G'_2$. Since G' is a proper subgraph G'' , we may assume without loss of generality that G_1 is a proper subgraph of G'_1 . But then G'_1 is G_1 -critical for $(8-g)$ -list-coloring. Hence by Theorem V.20, $d_{g,\varepsilon}(G'_1|G_1) \geq 0$. Let $G_3 = G'_1 \cup G_2$. Now $d_{g,\varepsilon}(G_3|G') = d_{g,\varepsilon}(G'_1|G_1) \geq 0$. By Proposition V.18, we have that

$$d_{g,\varepsilon}(G_3|H) = d_{g,\varepsilon}(G_3|G') + d_{g,\varepsilon}(G'|H) \geq 0 + 0 \geq 0,$$

contradicting the maximality of G' .

So we may assume that G'' is connected. By Corollary V.16, we have that $v(G'') \leq k'v(G') \leq k' \frac{g+\varepsilon}{\varepsilon}v(H)$ where k' is as in Corollary V.16. Since $H_1 \cup H_2 = H \subseteq G' \subseteq G''$ and G'' is connected, it follows that the distance between H_1 and H_2 is at most $v(G'') \leq k' \frac{5+\varepsilon}{\varepsilon}(v(H_1) + v(H_2)) = k(v(G_1) + v(G_2))$, as desired. \square

We are now ready to prove the strong hyperbolicity of \mathcal{F}_{345} as follows.

Proof of Theorem V.11. Since \mathcal{F}_{345} is hyperbolic by Theorem V.23, it suffices to prove the existence of a strong hyperbolic constant for \mathcal{F}_{345} . In fact, we prove the following stronger statement:

Claim V.25. *There exists $c' > 0$ such that the following holds: Let $G \in \mathcal{F}_{345}$ be a graph embedded in a surface Σ with rings \mathcal{R} . Let D_1, D_2 be two cycles of G such that there exists a cylinder $\Lambda \subseteq \Sigma$ with boundary components D_1 and D_2 . If G' is the subgraph of G consisting of all vertices and edges drawn in Λ , then*

$$v(G') \leq c'(v(D_1) + v(D_2)).$$

Proof. Specifically, we prove that $c' = \frac{5+\varepsilon}{\varepsilon}(k+1)$ suffices where ε is as in Theorem V.20 and k is as in Lemma V.24.

Now we regard G' as a planar graph. Let $H = D_1 \cup D_2$. Note that $v(H) \leq v(D_1) + v(D_2)$.

First suppose H is connected. Then by Lemma V.22, $d_{g,\varepsilon}(G'|H) \geq 0$. Thus by Lemma V.19, $v(G') \leq \frac{g+\varepsilon}{\varepsilon}v(H)$. But then

$$\begin{aligned} v(G') &\leq \frac{g+\varepsilon}{\varepsilon}v(H) \\ &\leq \frac{(5+\varepsilon)}{\varepsilon}(v(D_1) + v(D_2)) \leq c'(v(D_1) + v(D_2)), \end{aligned}$$

as desired.

So we may assume that H is not connected. First suppose that G' is not connected. It follows from Theorem V.1 that every component of G' contains a vertex in H' . Hence G' has exactly two components G_1 and G_2 . We may assume without loss of generality that G_1 contains D_1 and G_2 contains D_2 . By Lemma V.22, $d_{g,\varepsilon}(G_1|D_1) \geq 0$ and similarly $d_{g,\varepsilon}(G_2|D_2) \geq 0$. Thus by Theorem V.19, $v(G_1) \leq \frac{g+\varepsilon}{\varepsilon}v(D_1)$ and $v(G_2) \leq \frac{g+\varepsilon}{\varepsilon}v(D_2)$. But then

$$\begin{aligned} v(G') &\leq v(G_1) + v(G_2) \\ &\leq \frac{5+\varepsilon}{\varepsilon}(v(D_1) + v(D_2)) \leq c'(v(D_1) + v(D_2)), \end{aligned}$$

as desired.

So we may assume that G' is connected. Then by Lemma V.24, the distance between D_1 and D_2 is at most $k(v(D_1) + v(D_2)) \leq k(v(D_1) + v(D_2))$. Let P be a shortest path from D_1 to D_2 . Let $H' = H \cup P$. Now $v(H') \leq (k+1)(v(D_1) + v(D_2))$. Moreover, H' is connected. Thus by Lemma V.22, $d_{g,\varepsilon}(G'|H') \geq 0$. Hence by Lemma V.19, $v(G') \leq \frac{g+\varepsilon}{\varepsilon}v(H')$. But then

$$\begin{aligned} v(G') &\leq \frac{g+\varepsilon}{\varepsilon}v(H') \leq \frac{5+\varepsilon}{\varepsilon}(k+1)(v(D_1) + v(D_2)) \\ &= c'(v(D_1) + v(D_2)), \end{aligned}$$

as desired. \square

\square

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